

Webs, pockets, and buildings

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Based on joint work with subsets of *Christian Gaetz, Oliver Pechenik, Stephan Pfannerer, Jessica Striker, and Haihan Wu*

arXiv:2306.12501 (4-row)

arXiv:2402.13978 (2-column)

arXiv:2306.12506 (promotion permutations)

Slides: https://www.jpswanson.org/talks/2025_CanaDAM_pockets.pdf

Presented at the
Minisymposium on Web Graphs
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May 23rd, 2025

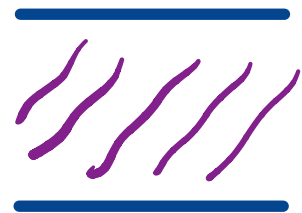
Outline

- SL_2 -webs and SL_3 -webs and bases
- Combinatorial geodesics
- SL_4 web basis and pockets
- Building embeddings

SL_2 -Webs

Def | An SL_2 -web is

- a bipartite graph embedded in



or

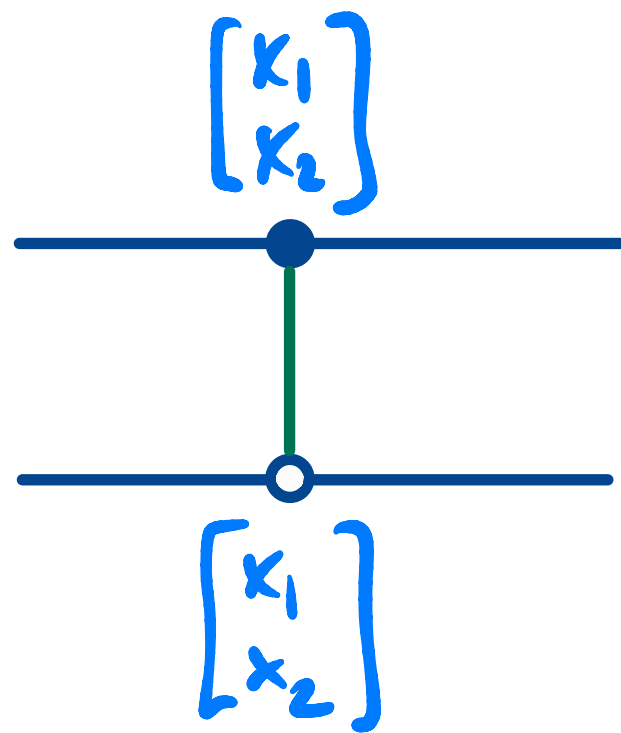


- with degree 2 internal vertices,
- and degree 1 boundary vertices.

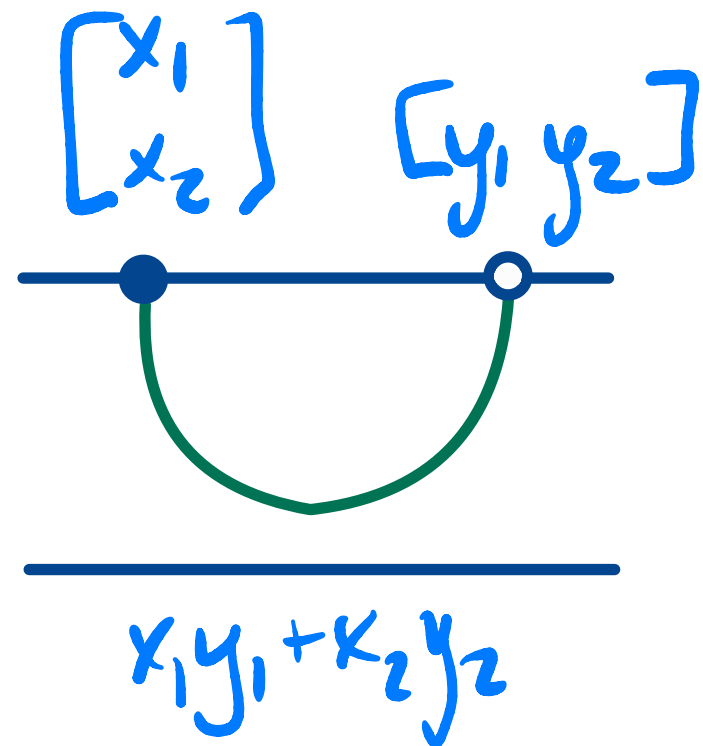
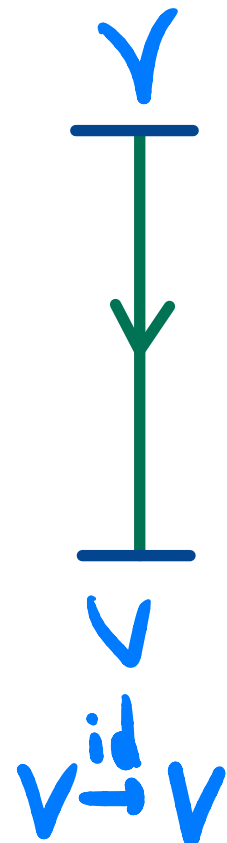
Fact | Fully encodes category of SL_2 -modules

SL_2 -Webs

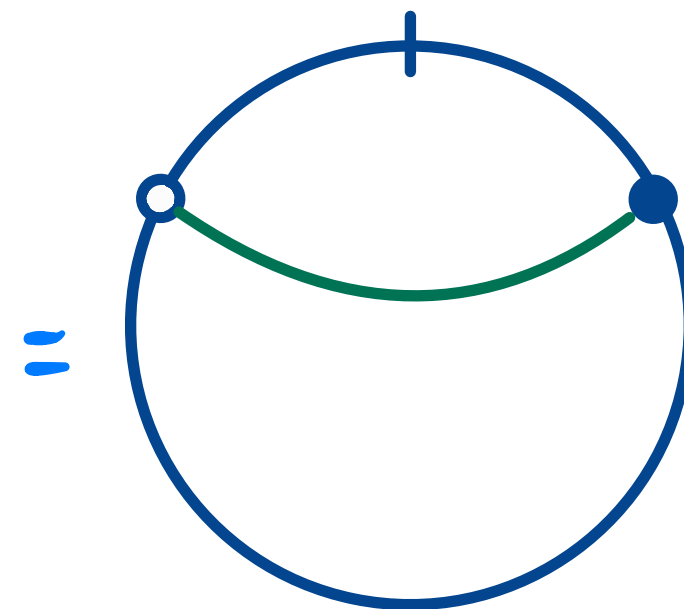
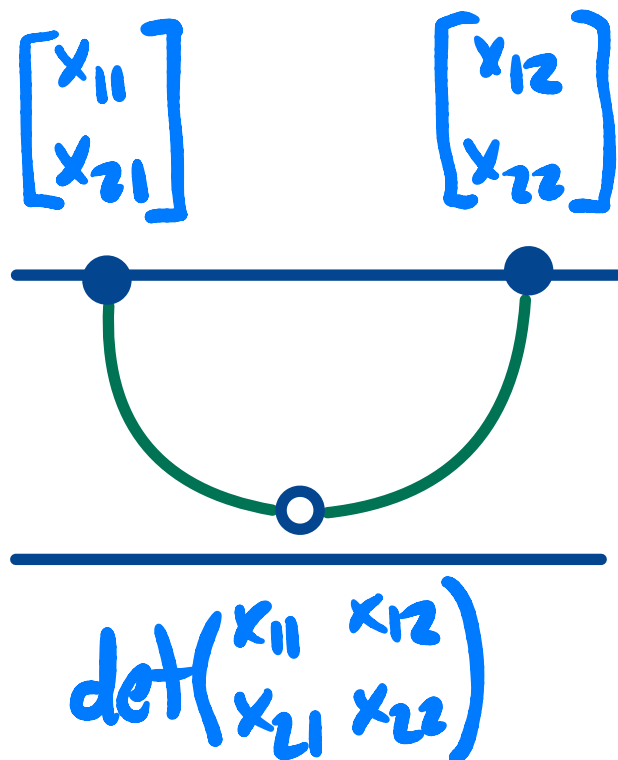
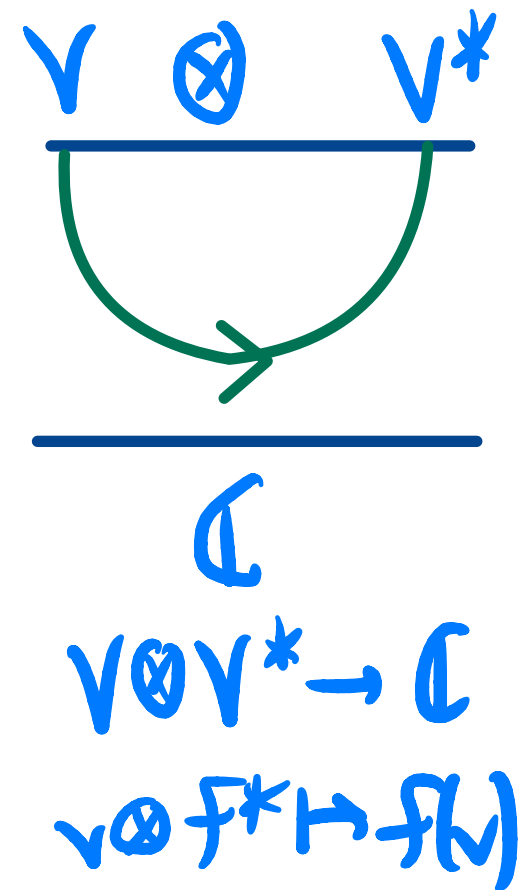
Ex Let $V = \mathbb{C}^2$.



=



=



SL_2 -Webs

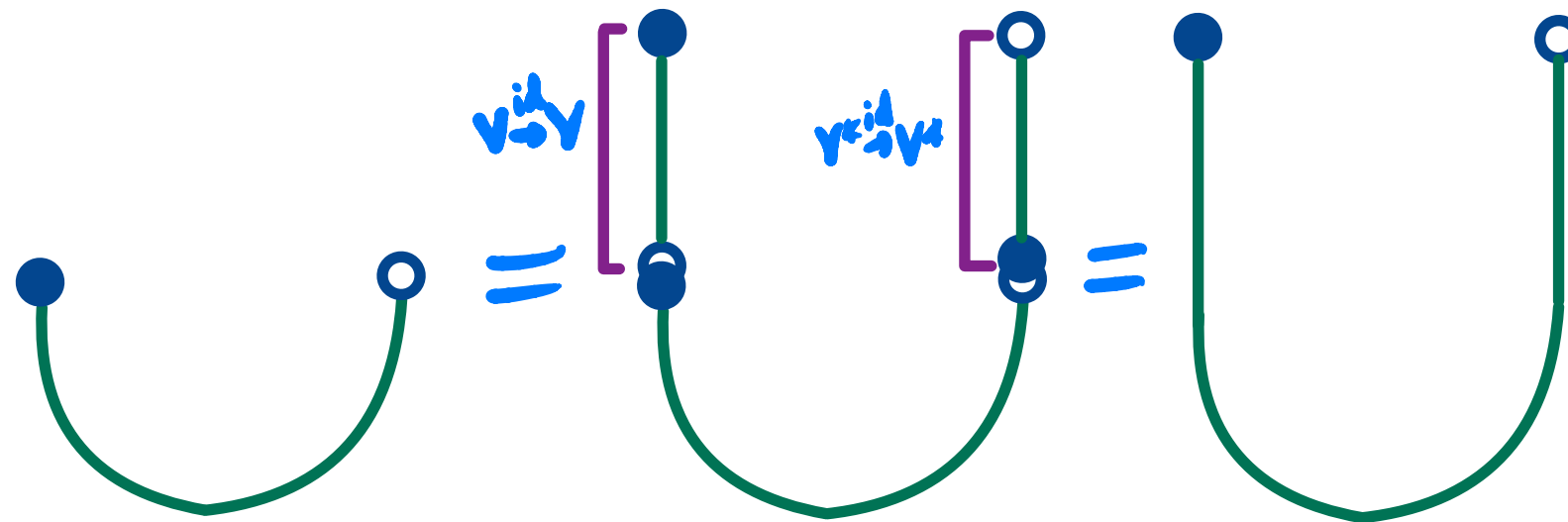
• Bipartite composition conventions

• $= V$ $\circ = V^*$ in domain

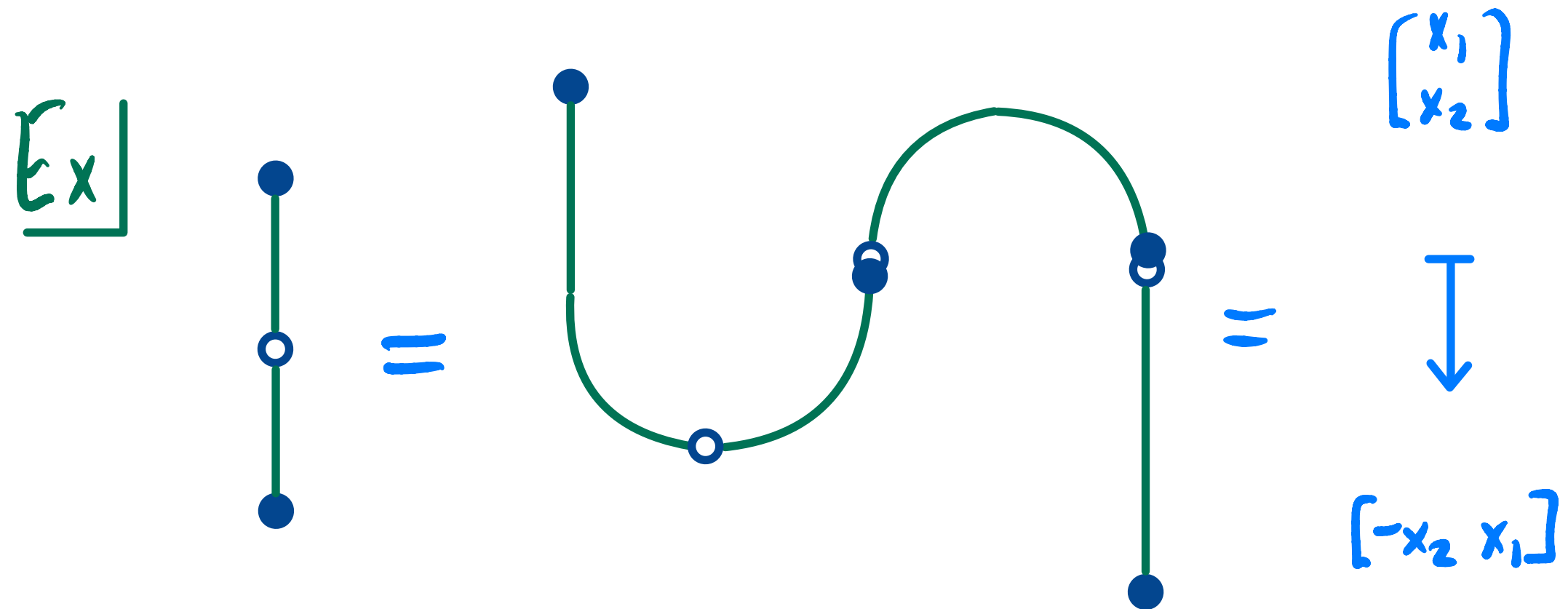
$\circ = V$ • $= V^*$ in codomain

— When composing, must match and cancel • pairs:

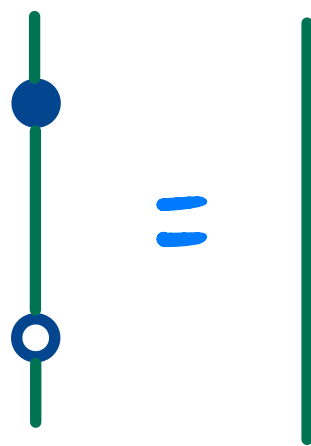
Ex



SL_2 -Webs



Ex Have contraction relation:



SL_2 -Webs

Facts | 1] Have well-defined map

$$\{SL(2) \text{ webs}\} / \text{isotopy} \longrightarrow \left\{ \begin{array}{c} SL(2) \text{ morphisms} \\ v^* \otimes v^* \otimes \dots \rightarrow v^* \otimes \dots \end{array} \right\}$$

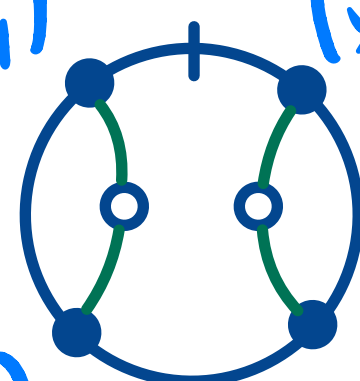
2] All of $\text{Rep}(SL_2)$ is encoded in disk case

(... surjective up to linear combinations, take Karoubi envelope...)

Q | Generators and relations? Bases??

SL_2 -Webs

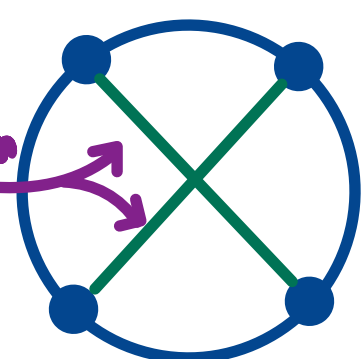
Ex

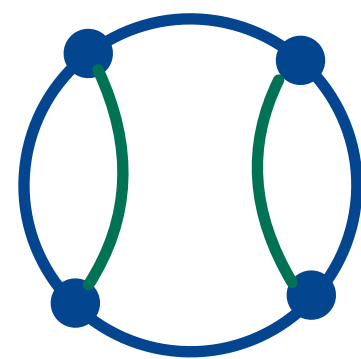


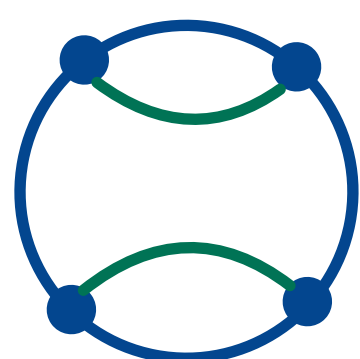
$$= \det \begin{pmatrix} x_{11} & x_{12} \\ x_{21} & x_{22} \end{pmatrix} \cdot \det \begin{pmatrix} x_{13} & x_{14} \\ x_{23} & x_{24} \end{pmatrix} \in \text{Inv}(V^{\otimes 4})$$

$$= \text{Hom}_{SL_2}(V^{\otimes 4}, \mathbb{C})$$

Ex Plücker relations:



$$=$$


$$+$$


$$(x_{11}x_{23} - x_{21}x_{13})(x_{12}x_{24} - x_{12}x_{14})$$

$$=$$

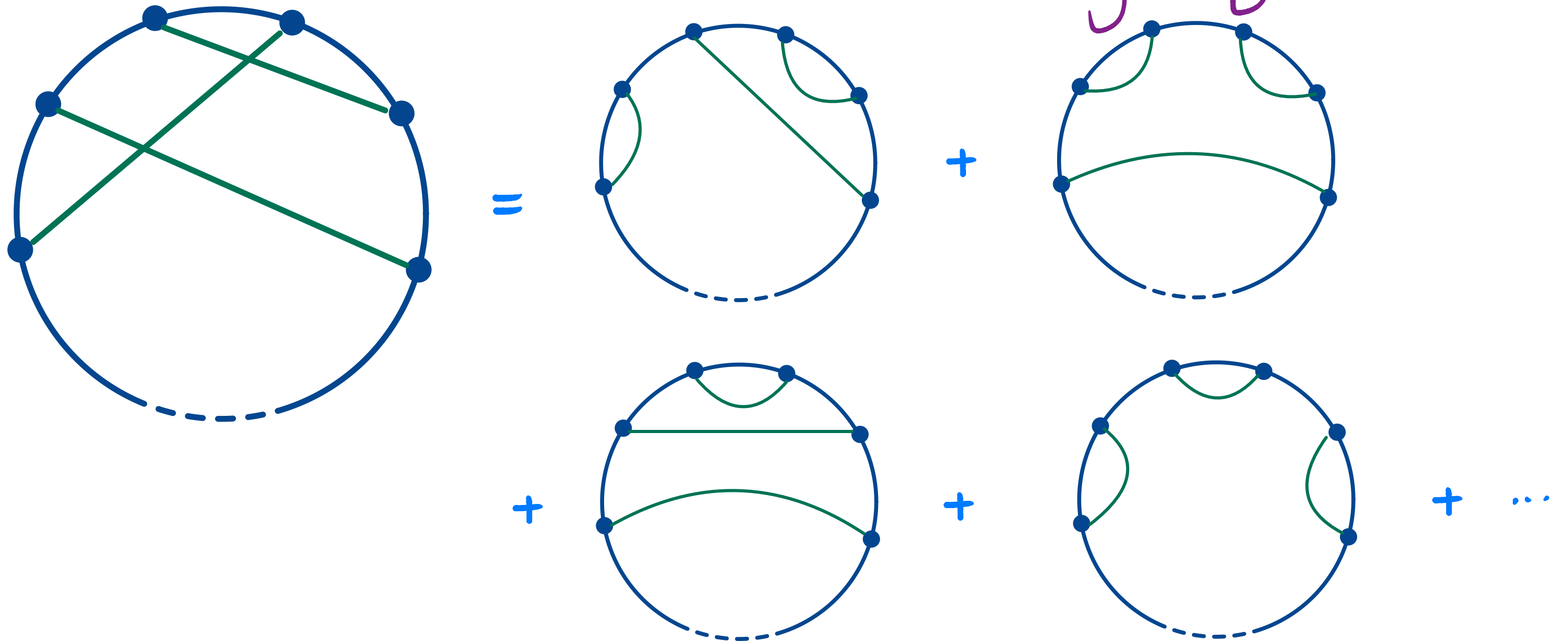
$$(x_{11}x_{22} - x_{21}x_{12})(x_{13}x_{24} - x_{23}x_{14})$$

$$+$$

$$(x_{11}x_{24} - x_{21}x_{14})(x_{12}x_{23} - x_{22}x_{13})$$

Temperley-Lieb basis

- Using $\text{diagram with crossing} = \text{diagram with two arcs} + \text{diagram with two arcs}$, can reduce any matching diagram to a linear combination of matching diagrams:



Temperley-Lieb basis

Thm The noncrossing 2-row webs are a basis for
 $\text{Inv}_{\text{SL}_2}(V_1 \otimes \dots \otimes V_n)$ $(V_i \in \{V, V^*\})$

called the Temperley-Lieb basis.

PF • Spanning: diagrams span by classical invariant theory,
noncrossing by uncrossing rule.

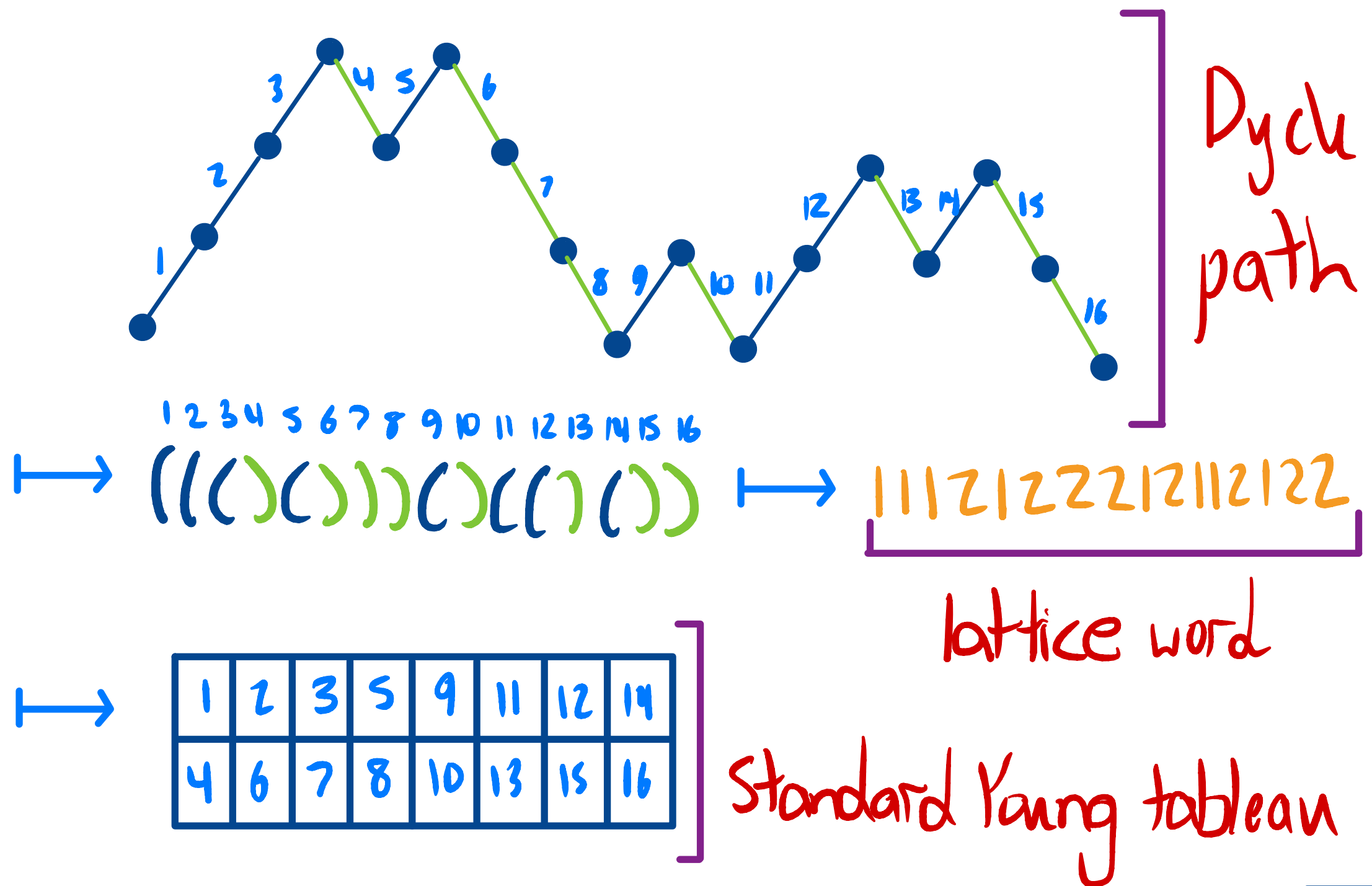
• Independence: by Pieri rule,

$$\dim \text{Inv}_{\text{SL}_2}(V^{\otimes n}) = \# \text{ST}(2 \times \frac{n}{2}). \text{ Count!}$$

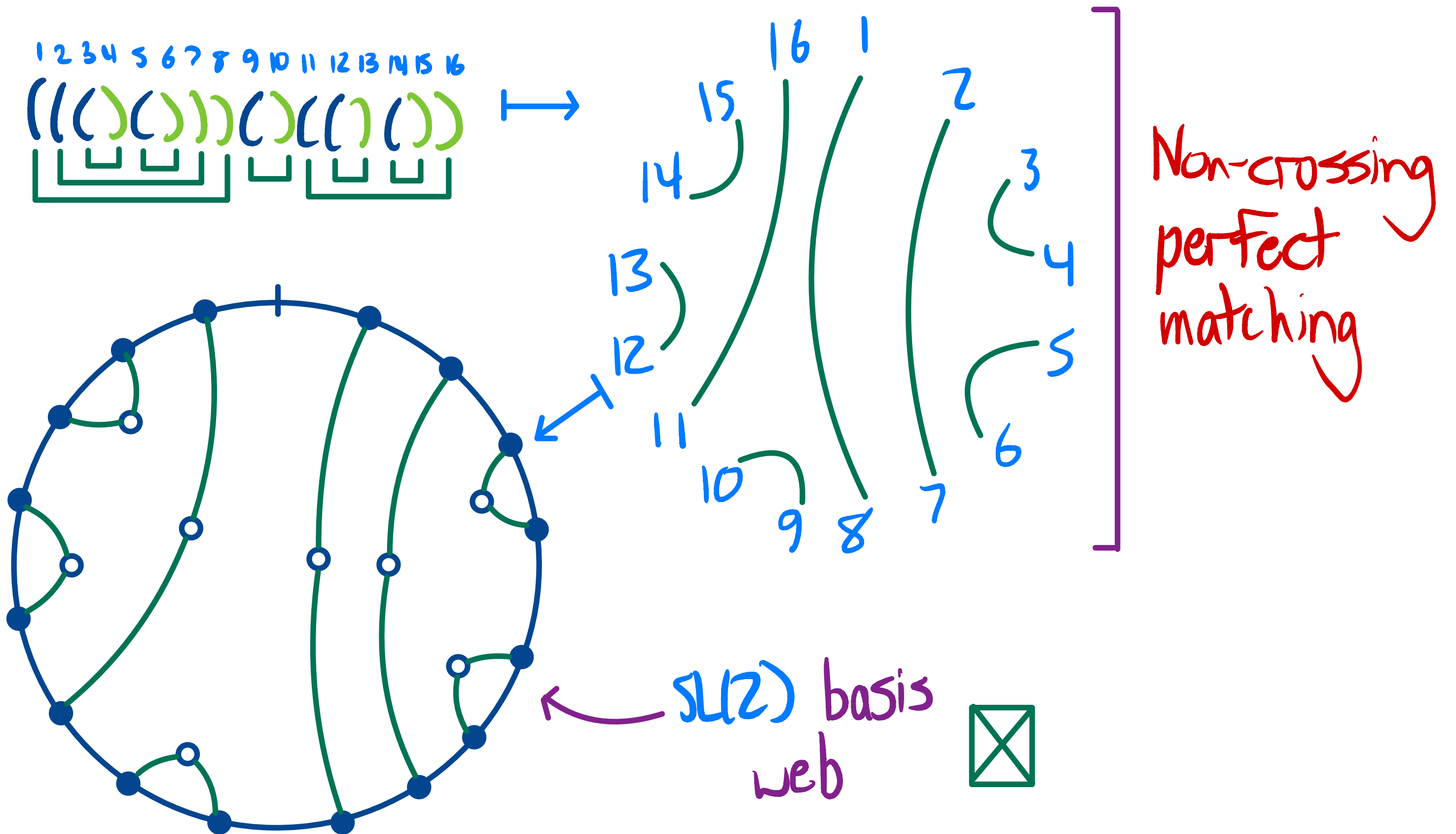


Temperley-Lieb basis

Some Catalan bijections:

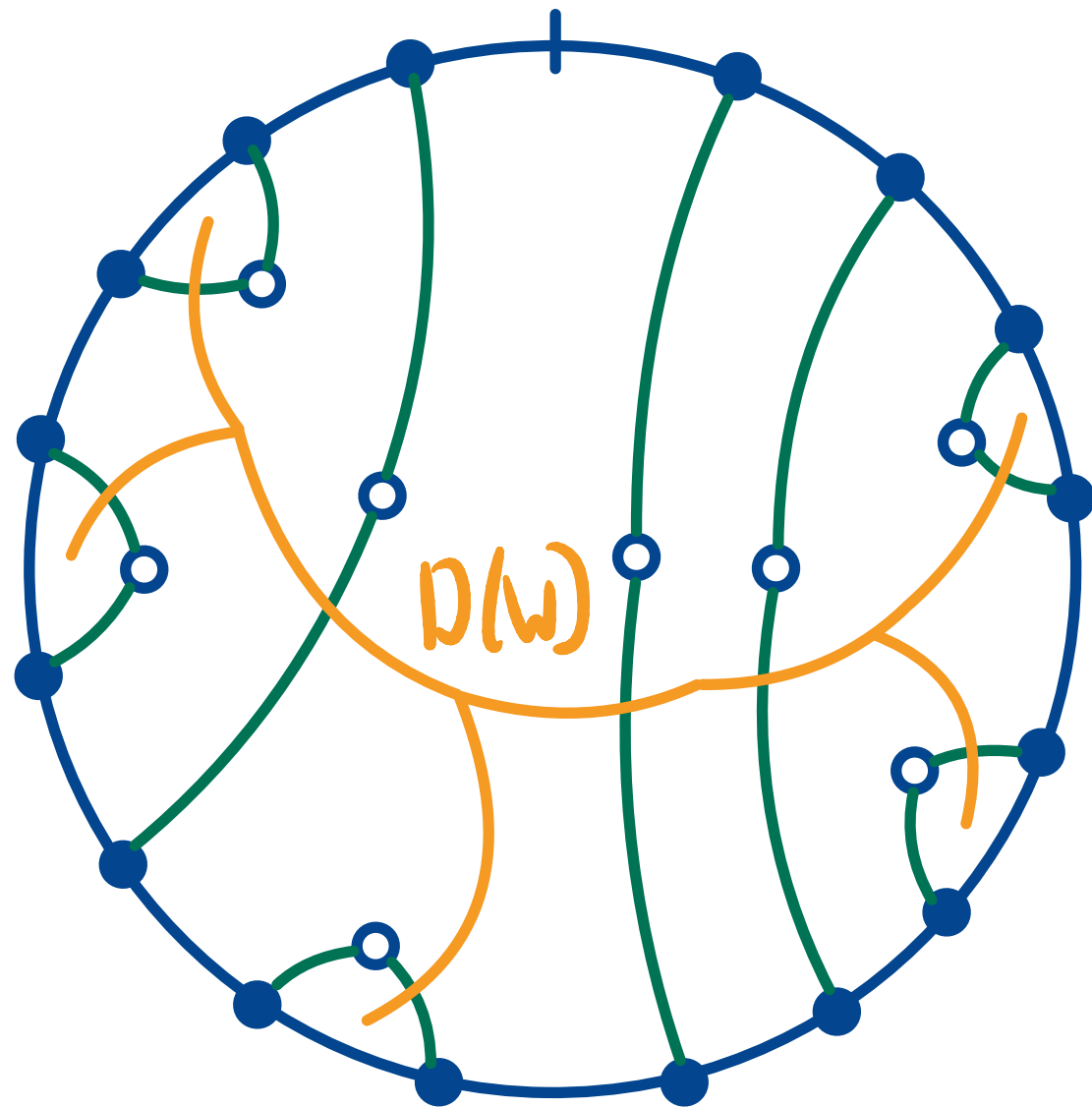


Temperley-Lieb basis



Duals and trees

Obs | The dual graph $D(W)$ of an $SL(2)$ basis web W is a tree:



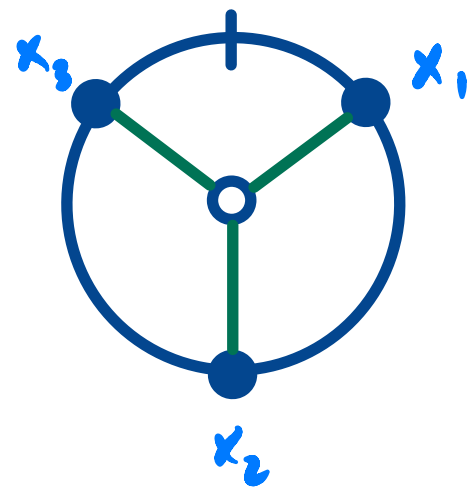
e.g. #faces of W
= #vertices of $D(W)$

SL_3 -webs

- Let $V = \mathbb{C}^3$, $V_i \in \{V, V^*\}$.

Q Is there a nice web basis for $\text{Inv}_{SL_3}(V_1 \otimes \dots \otimes V_n)$?

- Use bipartite planar graphs in a disk built from



$$= \det \begin{pmatrix} | & | & | \\ x_1 & x_2 & x_3 \\ | & | & | \end{pmatrix} \text{ (so trivalent)}.$$

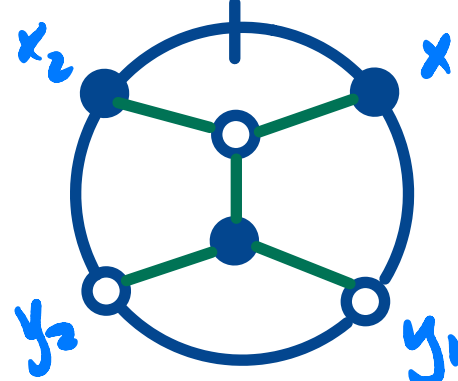
(+ dual)

Have univalent boundary vertices
and connected to boundary.

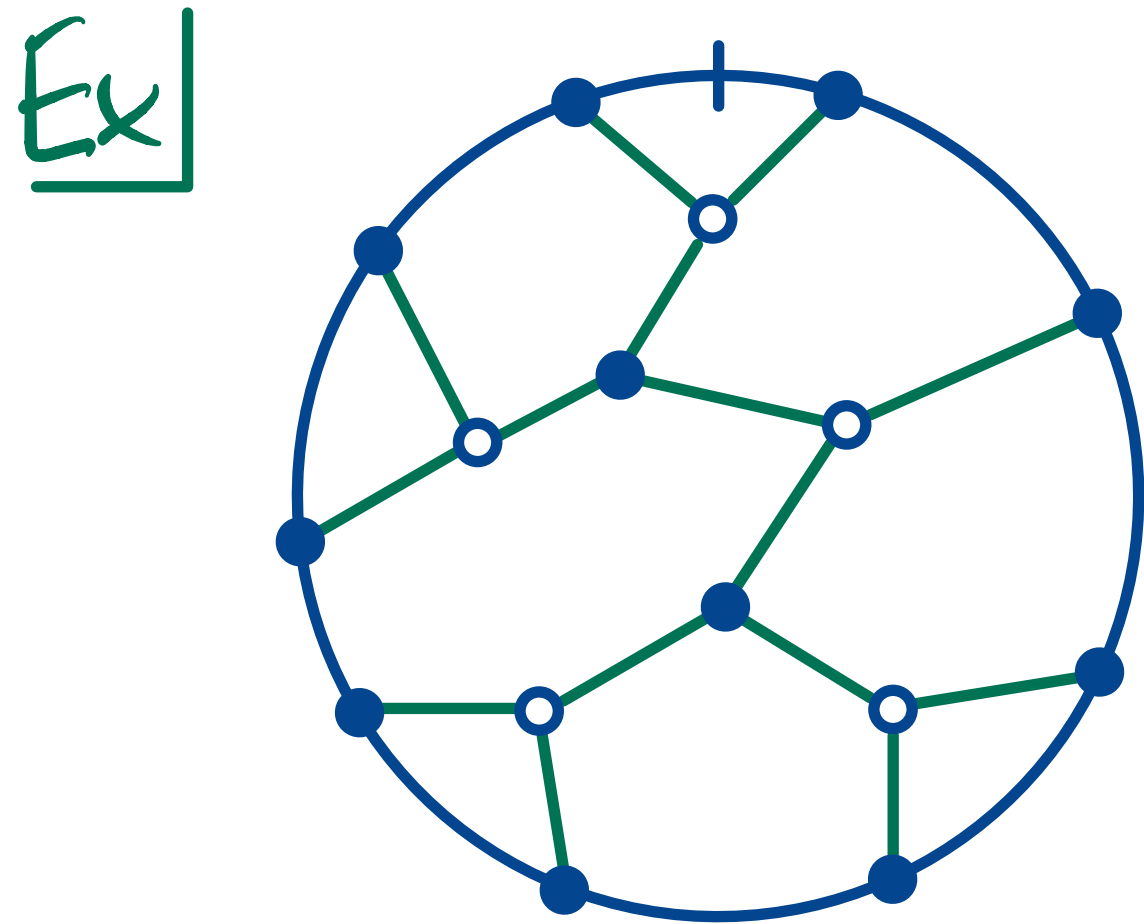
SL_3 -webs

SL_3 -webs

Ex



$$= \det \begin{pmatrix} 1 & 1 & 1 \\ x_1 & x_2 & y_1 \wedge y_2 \\ 1 & 1 & 1 \end{pmatrix}$$



= ... a polynomial obtained by
summing over proper edge
3-colorings...

SL_3 -webs

Thm (Kuperberg '94) The generating SL_3 -web relations are

$$\bigcirc = 3$$

$$\text{---} \bigcirc \text{---} = 2 \cdot \text{---}$$

$$\begin{array}{c} \diagup \quad \diagdown \\ \bigcirc \quad \bullet \\ | \quad | \\ \bullet \quad \bigcirc \\ \diagdown \quad \diagup \end{array} = \begin{array}{c} \diagup \quad \diagdown \\ \text{---} \quad \text{---} \\ \diagdown \quad \diagup \end{array} + \begin{array}{c} \diagup \quad \diagdown \\ \text{---} \quad \text{---} \\ \diagdown \quad \diagup \end{array}$$

Non-elliptic web basis

Thm (Kuperberg '94)

Call an SL_3 -web nonelliptic if it has
no 2-faces or 4-faces.

The nonelliptic webs form a basis of

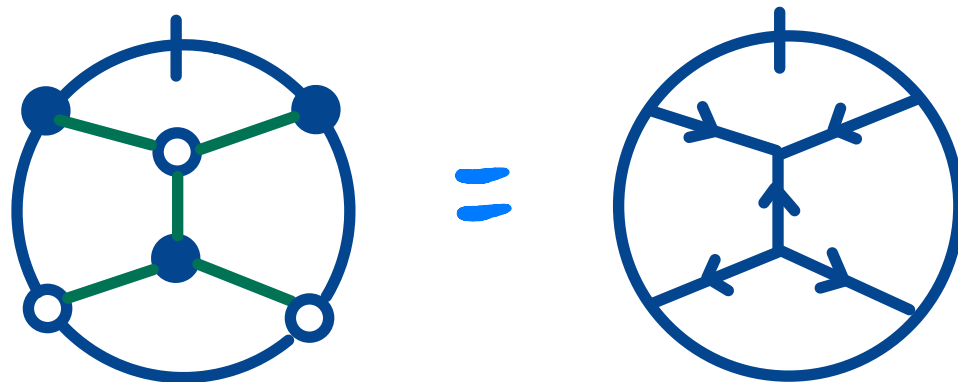
$$\text{Inv}_{SL_3}(V_1 \otimes \dots \otimes V_n). \quad (v_i \in \{V, V^*\})$$

Non-elliptic web basis

PF • Spanning: similar to SL_2 case.

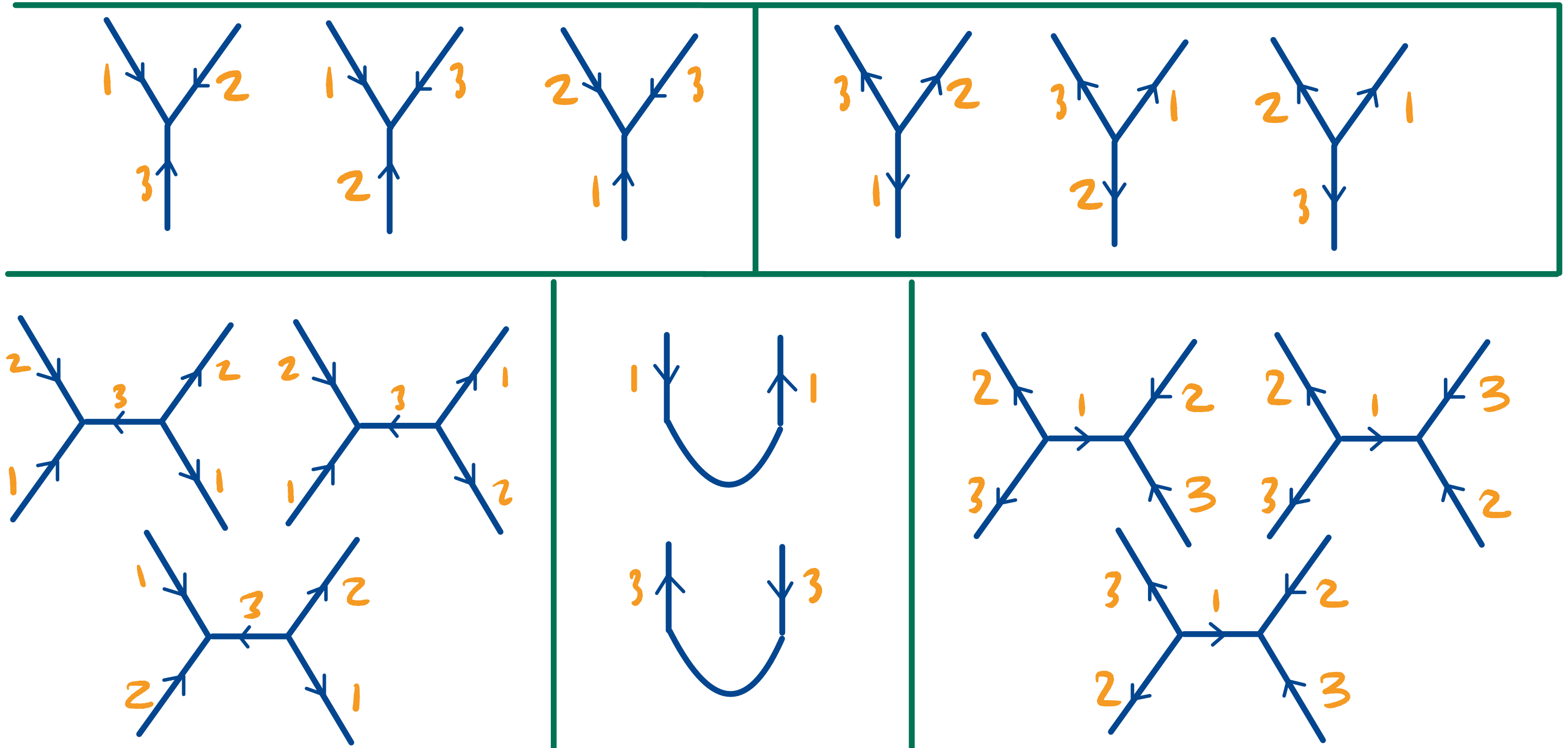
• Independence: bijection to $SVT(3 \times \frac{1}{3})$ using growth rules. ($\exists$ other approaches)

Will use directed notation here:



SL_3 -growth rules

Kuperberg-Khoranav growth rules:



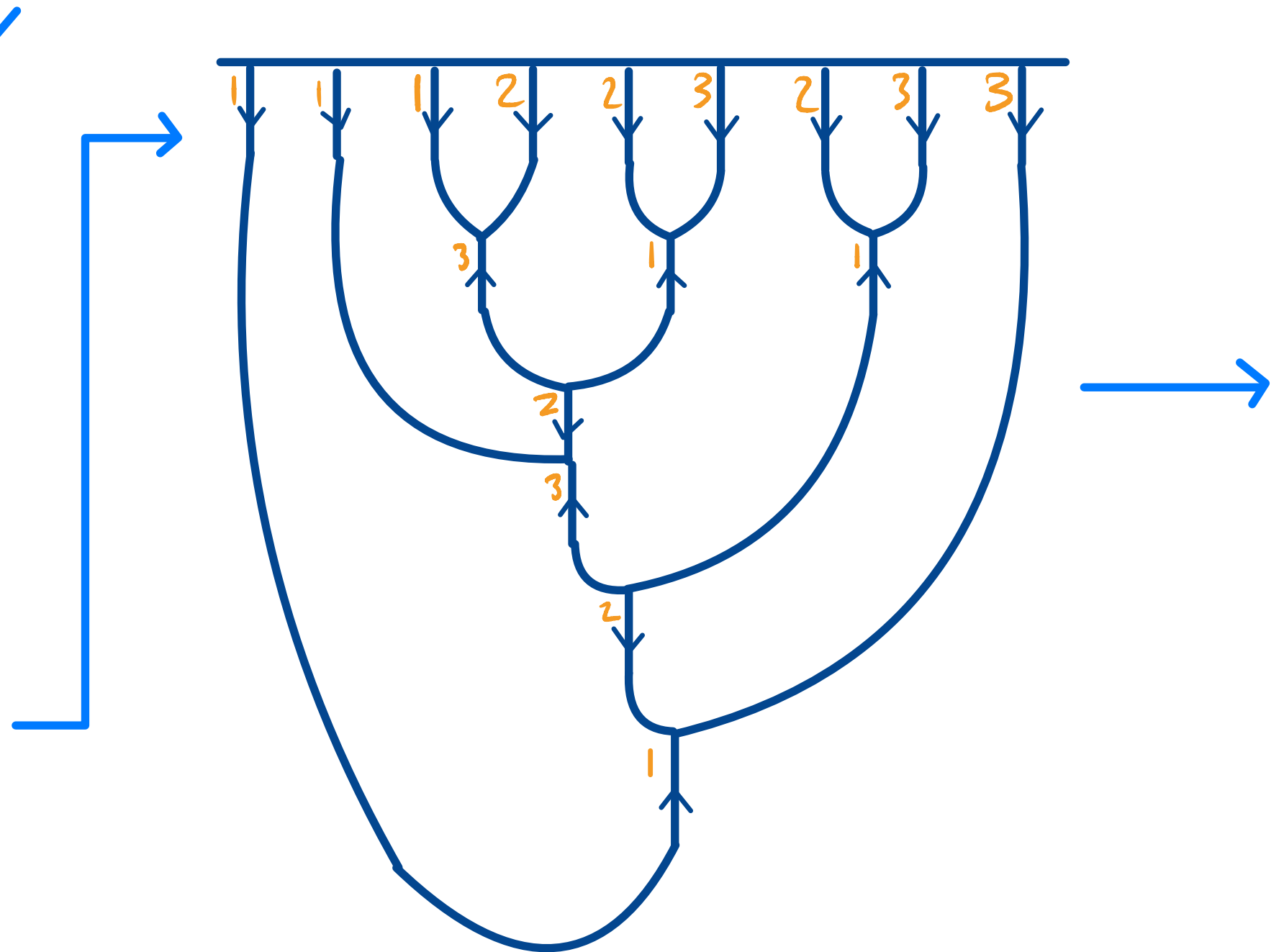
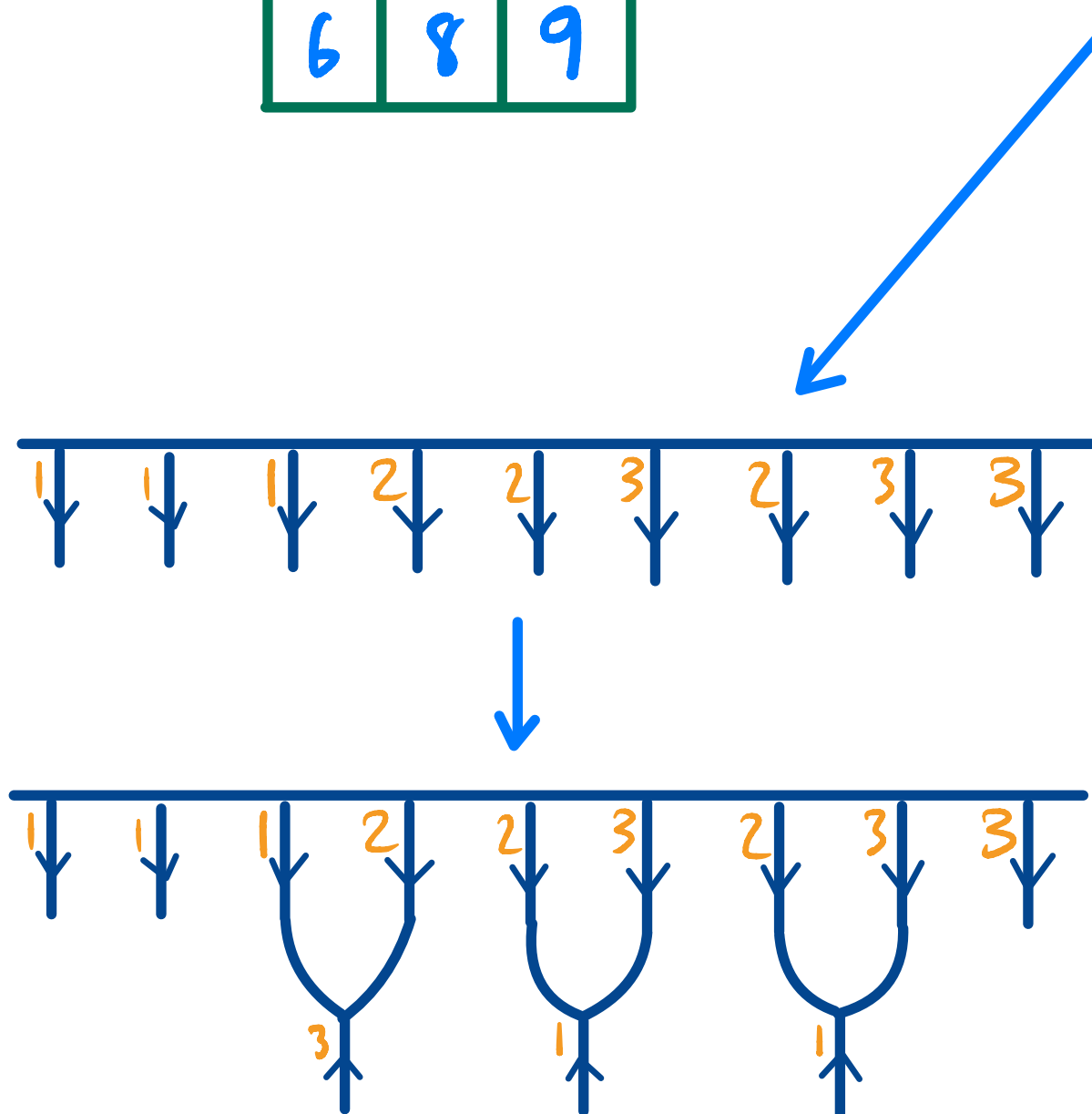
SL_3 -growth rules

Ex

$T =$

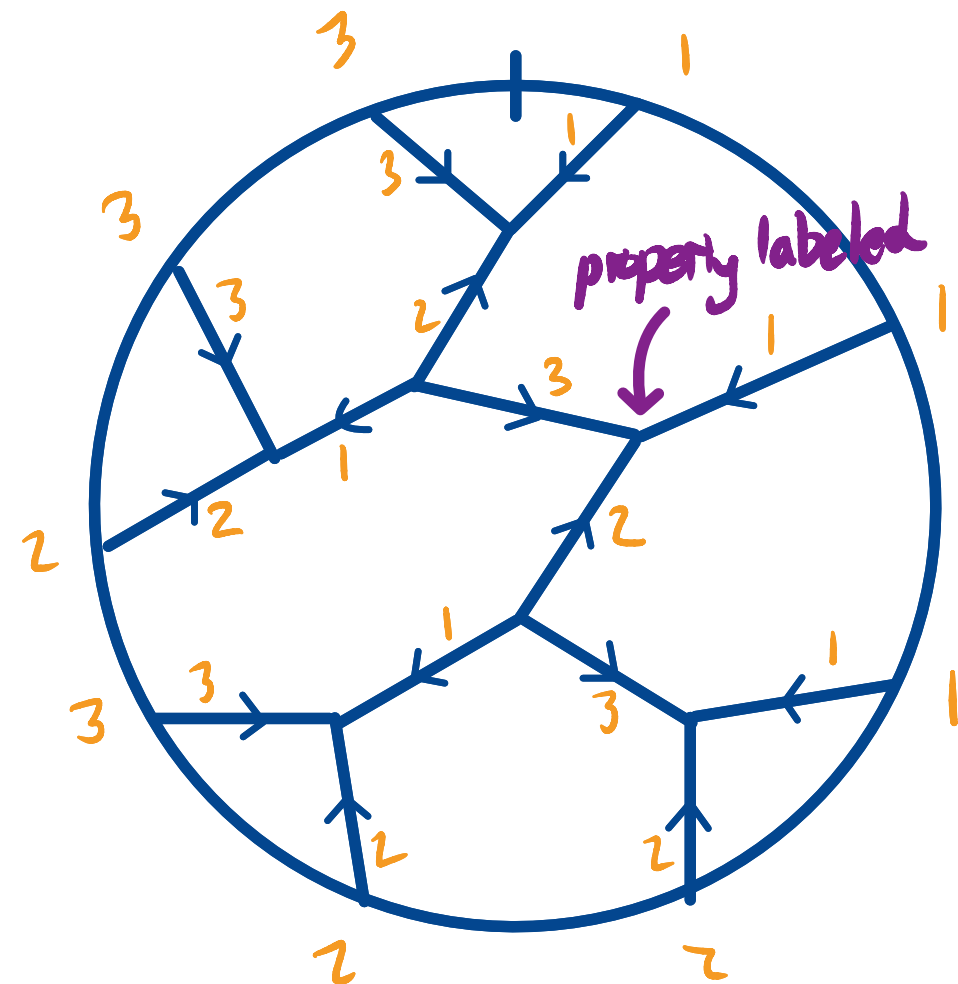
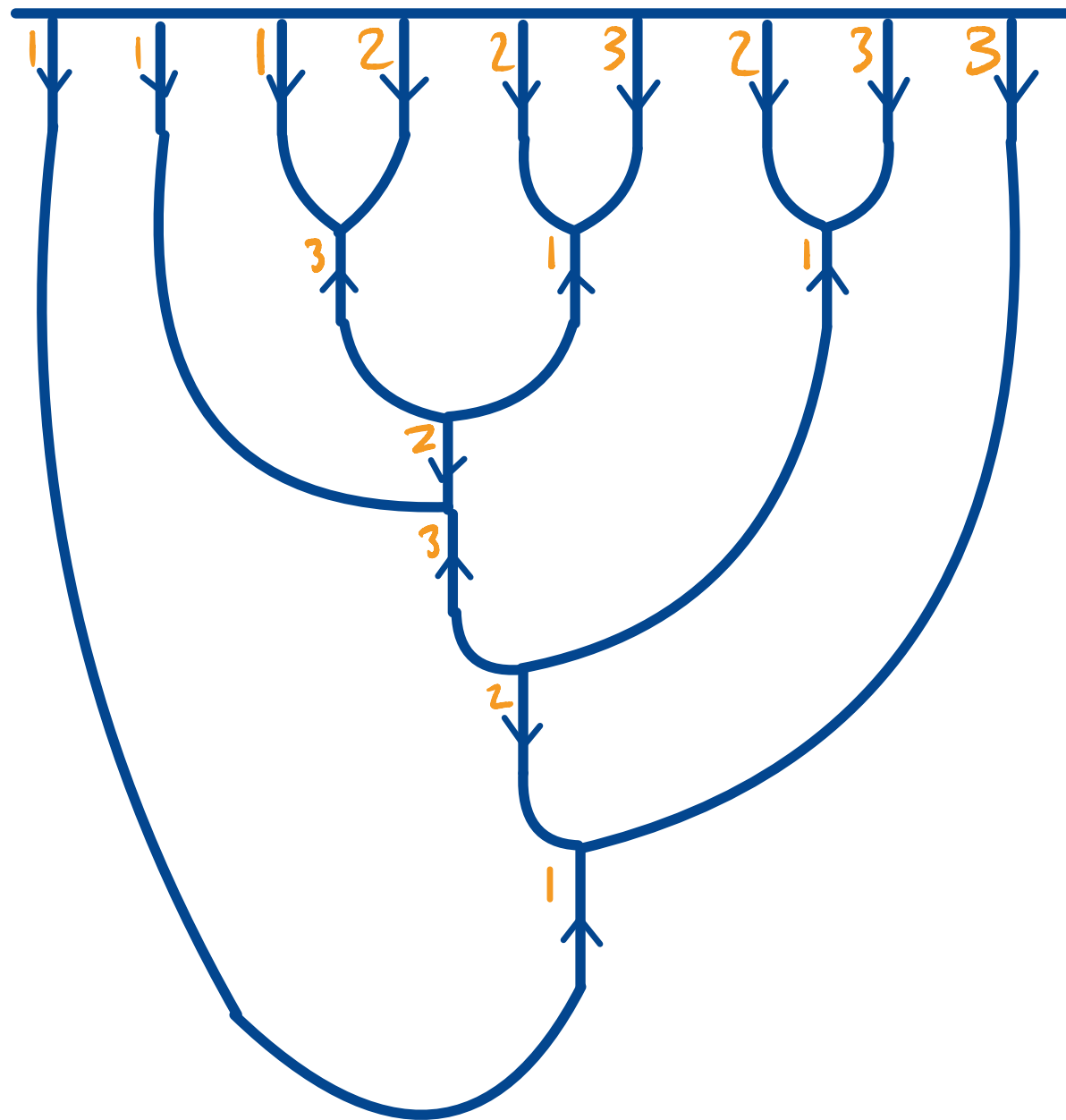
1	2	3
4	5	7
6	8	9

$\rightarrow 111223233$



SL_3 -growth rules

($T \rightarrow 111223233$)



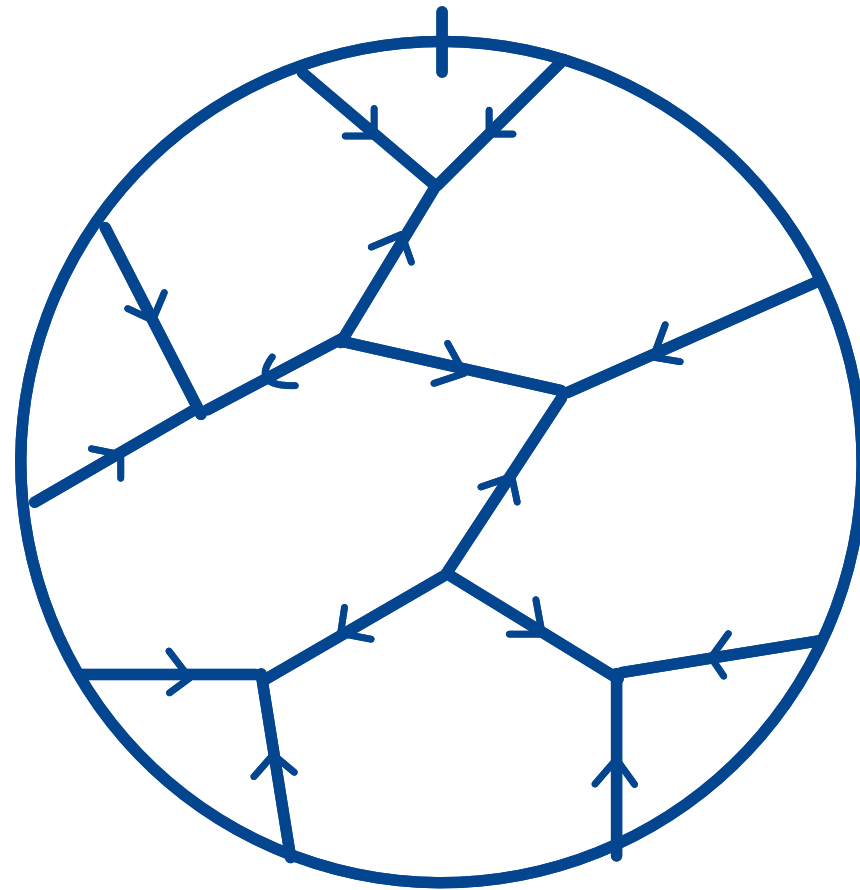
Now just erase labels!

SL_3 -growth rules

In all:

$T =$

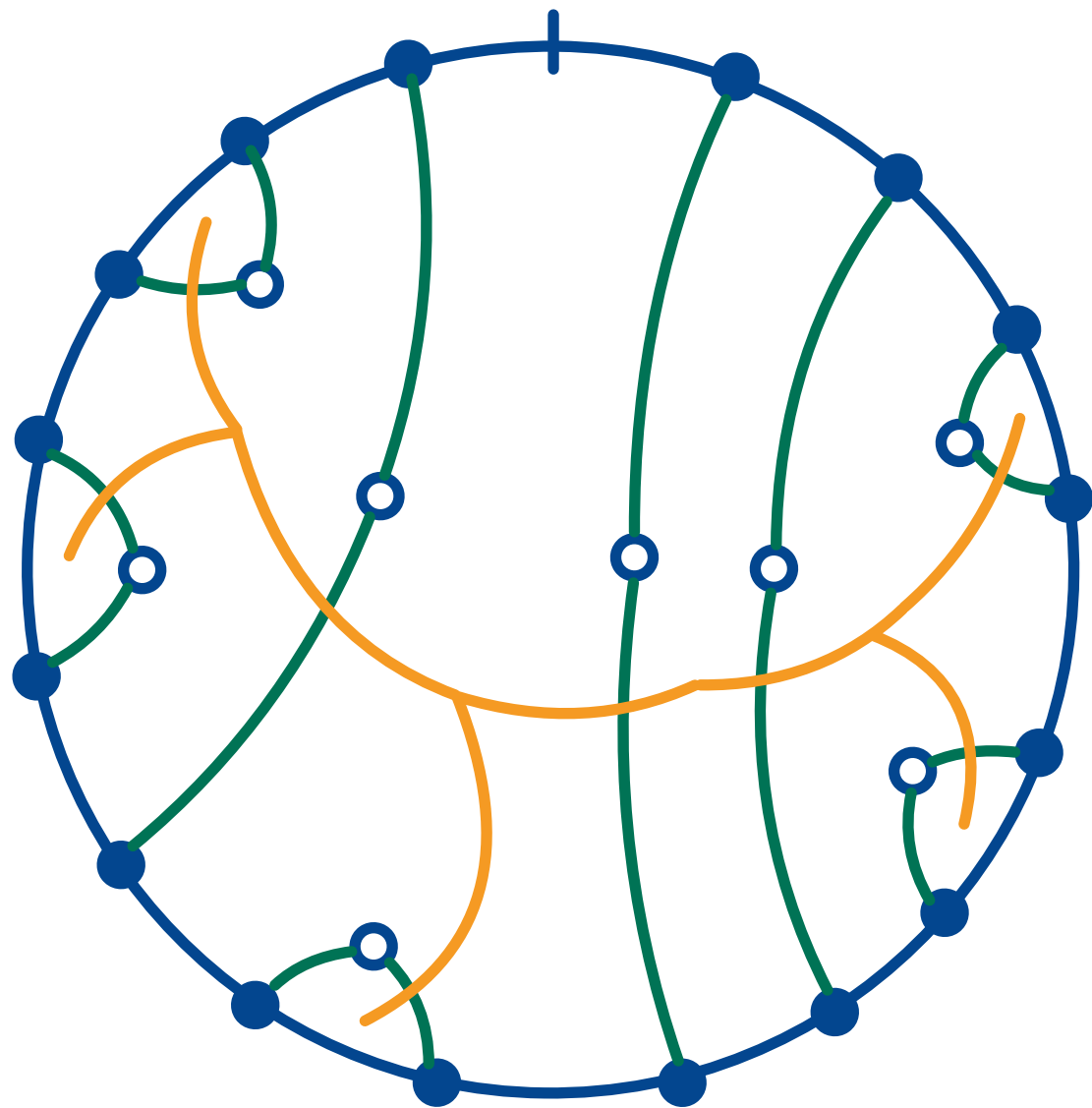
1	2	3
4	5	7
6	8	9



Q What did the labels mean?

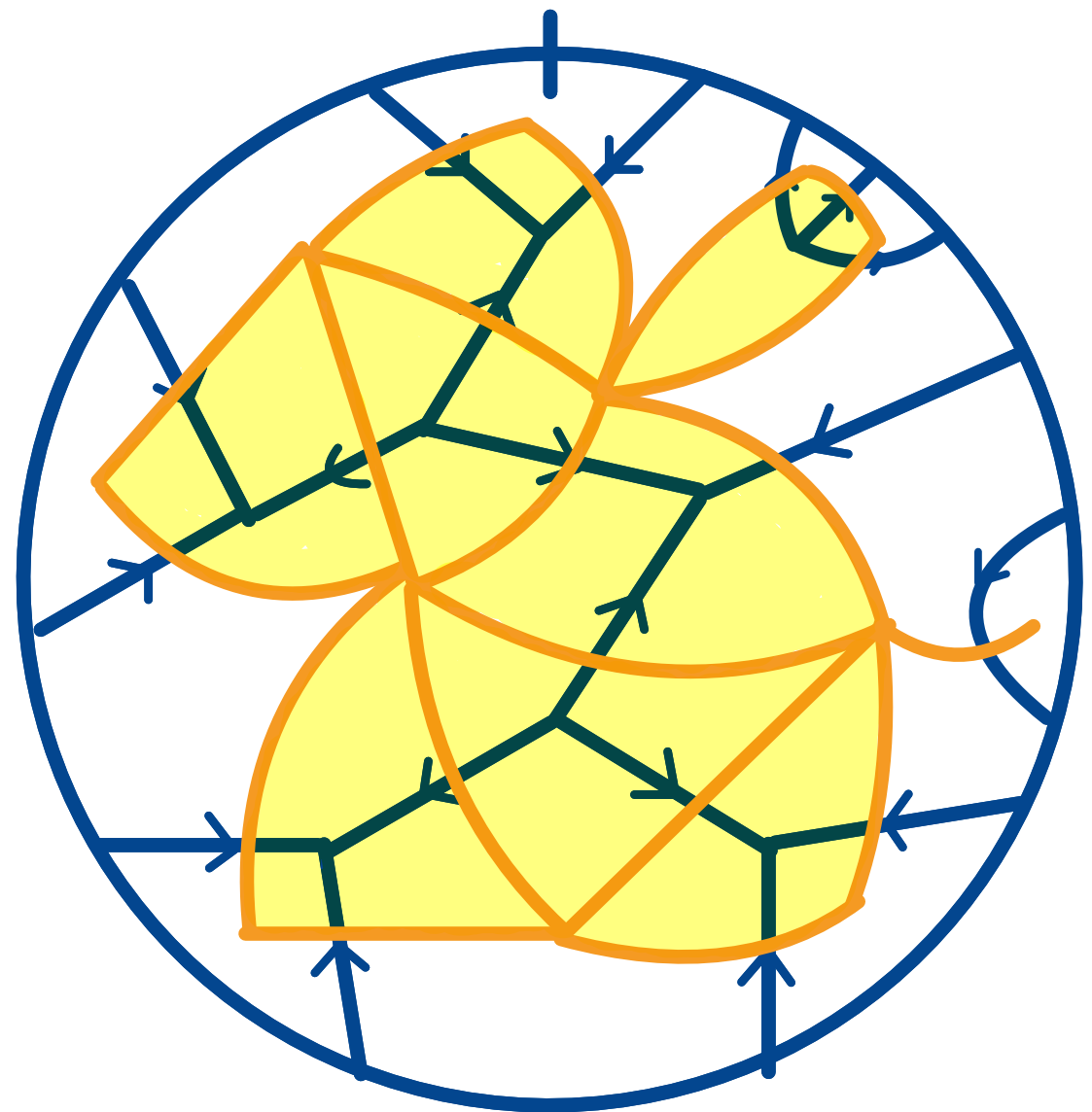
SL_3 -web duals

Obs | $D(W)$ for $SL(3)$ is a triangulation:



$D(W)$ is 1D for $SL(2)$

vs.



$D(W)$ is "2D" for $SL(3)$

SL_3 -web duals

Next

• Fontaine-Kannitzer-Kuperberg '13 showed duals of

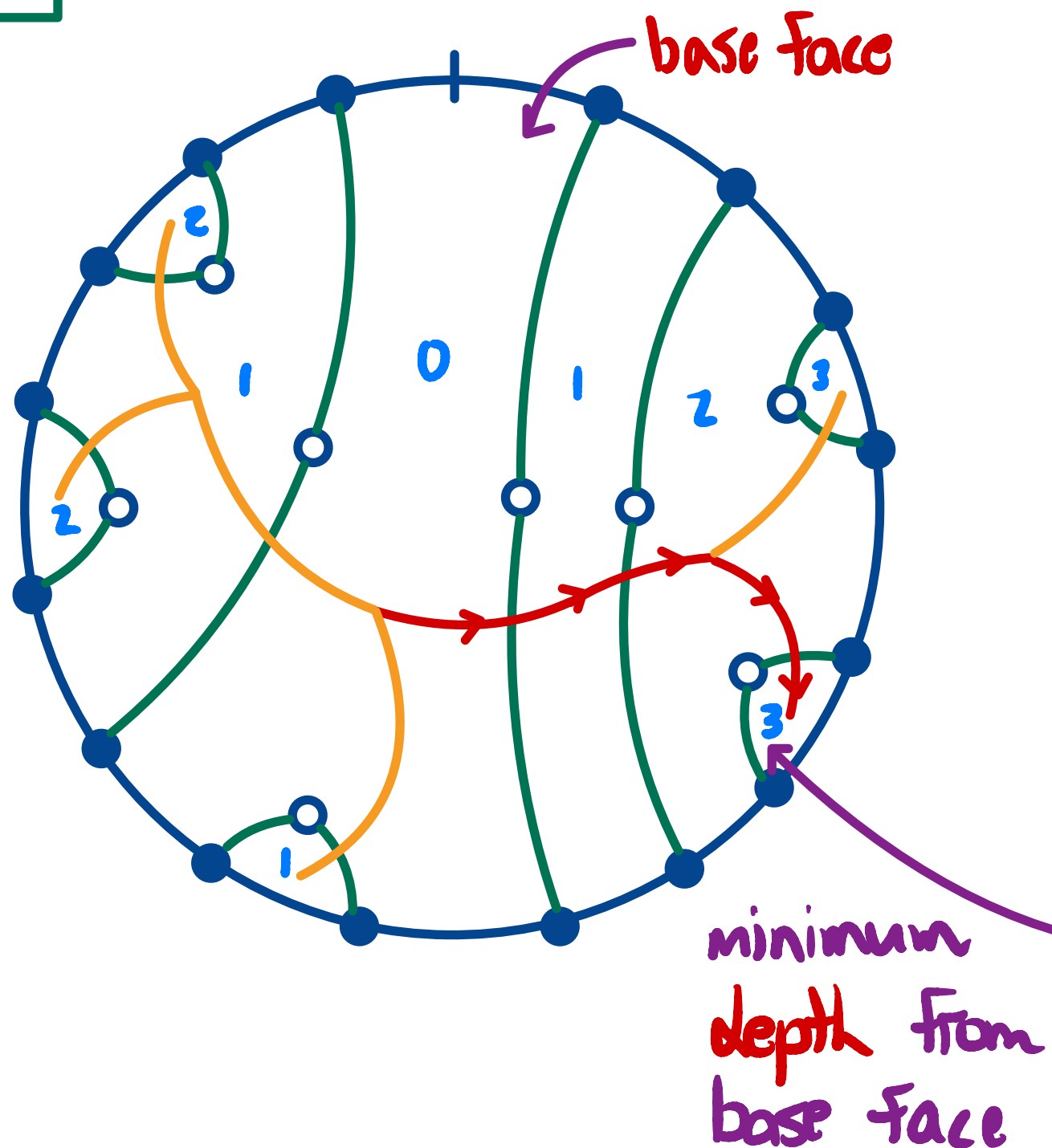
SL_3 basis webs live inside the

affine building $\Delta(SL_3^\vee)$

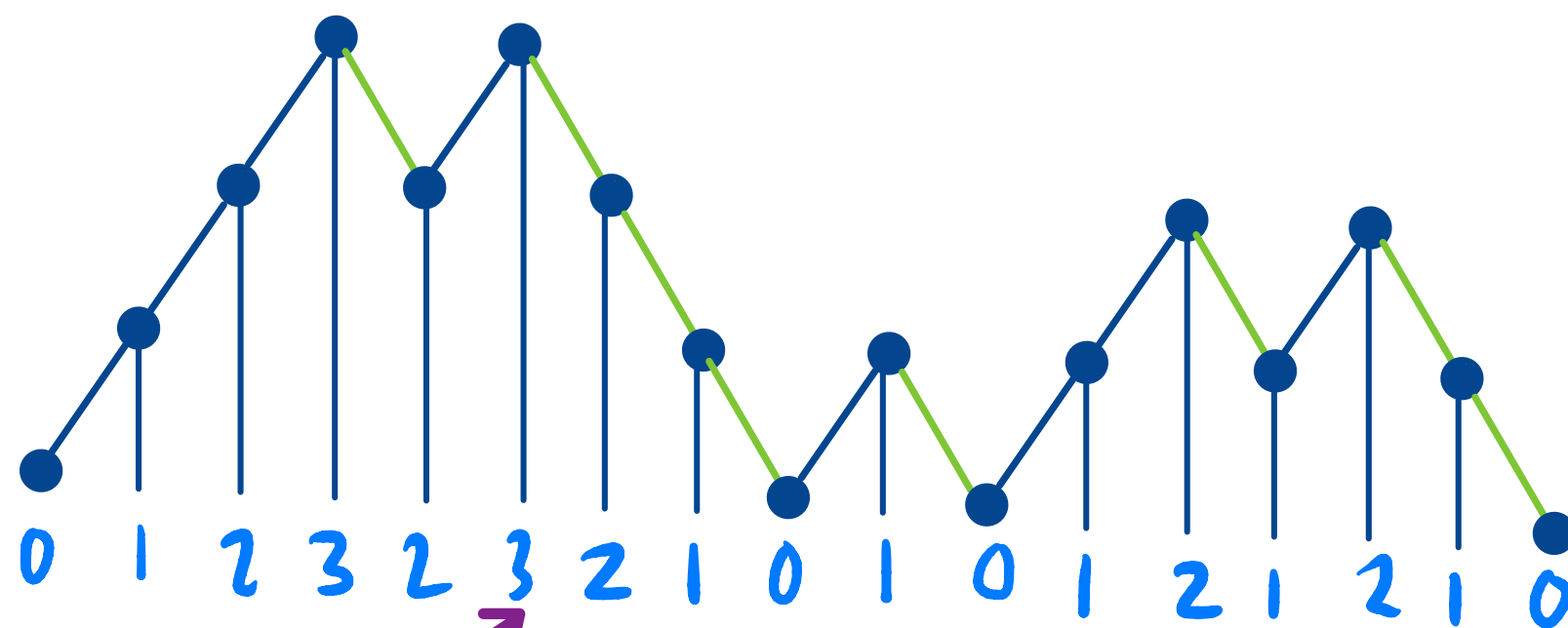
- Growth rules build one triangle at a time
- Labels encode distance information
- Non-elliptic condition $\Leftrightarrow AT(0)$

Distances

E_x



1	2	3	5	9	11	12	14
4	6	7	8	10	13	15	16



- Hence depths around the boundary encode the tableau

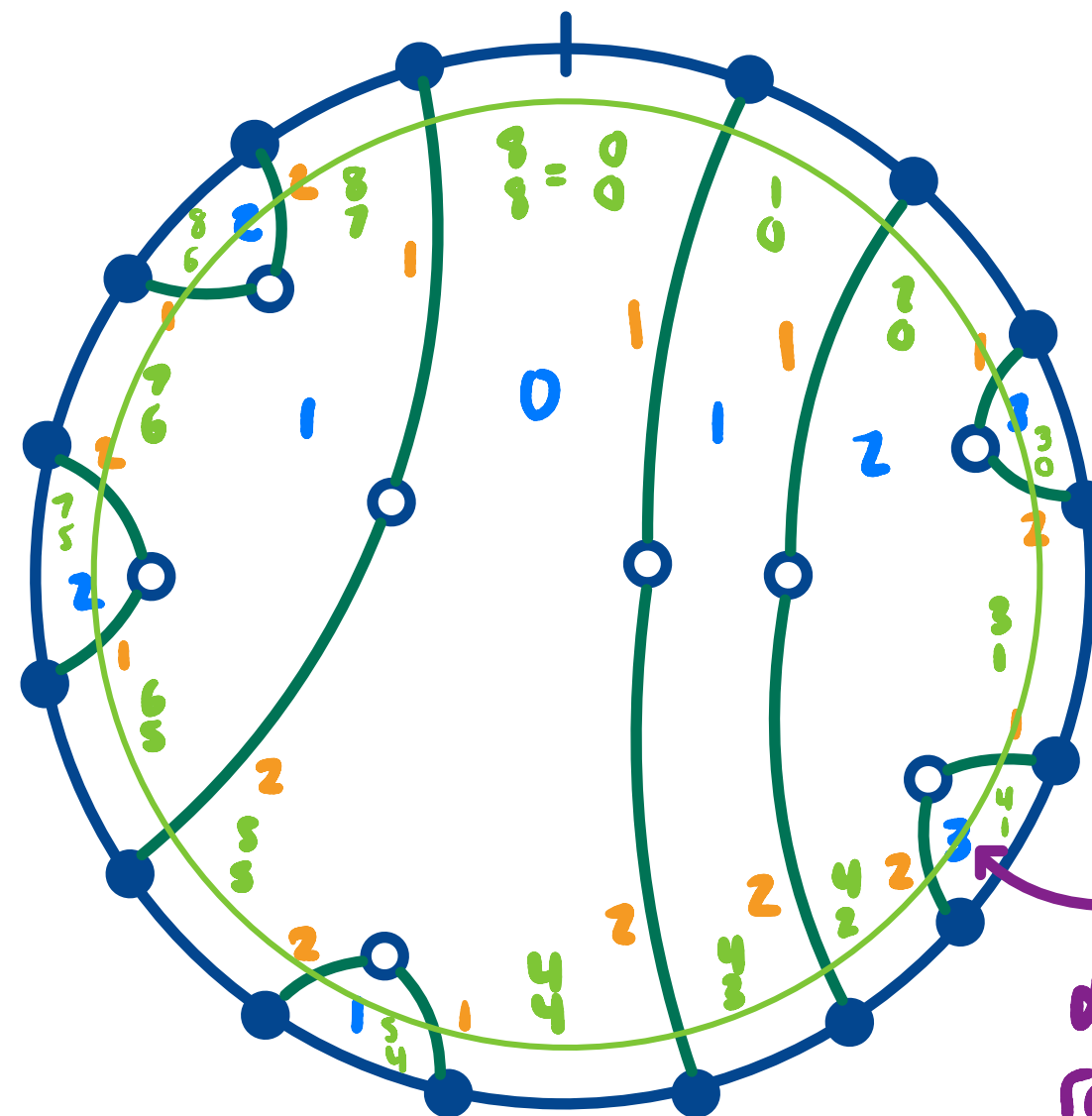
Distances

Ex

1	2	3	5	9	11	12	14
4	6	7	8	10	13	15	16



$\phi \subset (1,0) \subset (2,0) \subset (3,0) \subset (3,1)$
 $\subset \dots \subset (8,8)$



Label so that boundary is
 1112122212112122

Distances

Recall | The $SL(r)$ dominant weights are
 $\lambda = (\lambda_1 \geq \dots \geq \lambda_r) \in \mathbb{Z}^r / \langle (1, \dots, 1) \rangle$.

• The fundamental weights of $SL(r)$ are

$$\omega_i = (1^i, 0^{r-i})$$

$$\omega_i^* = (0^{r-i}, -1^i) \equiv \omega_{r-i}.$$

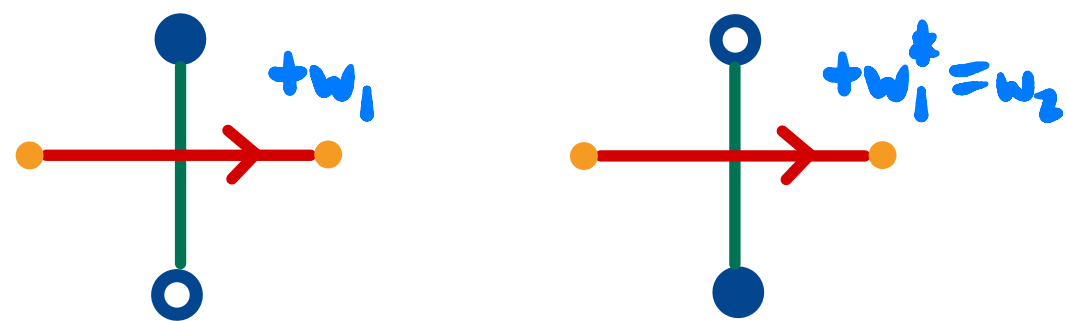
• If $\lambda, \mu \in \mathbb{Z}_{\text{dom}}^r$ where $|\lambda| = |\mu| + rk$ for some $k \in \mathbb{Z}$,

say

$$\lambda \geq \mu \Leftrightarrow \forall m, \sum_{j=1}^m \lambda_j \geq \sum_{j=1}^m (\mu_j + k).$$

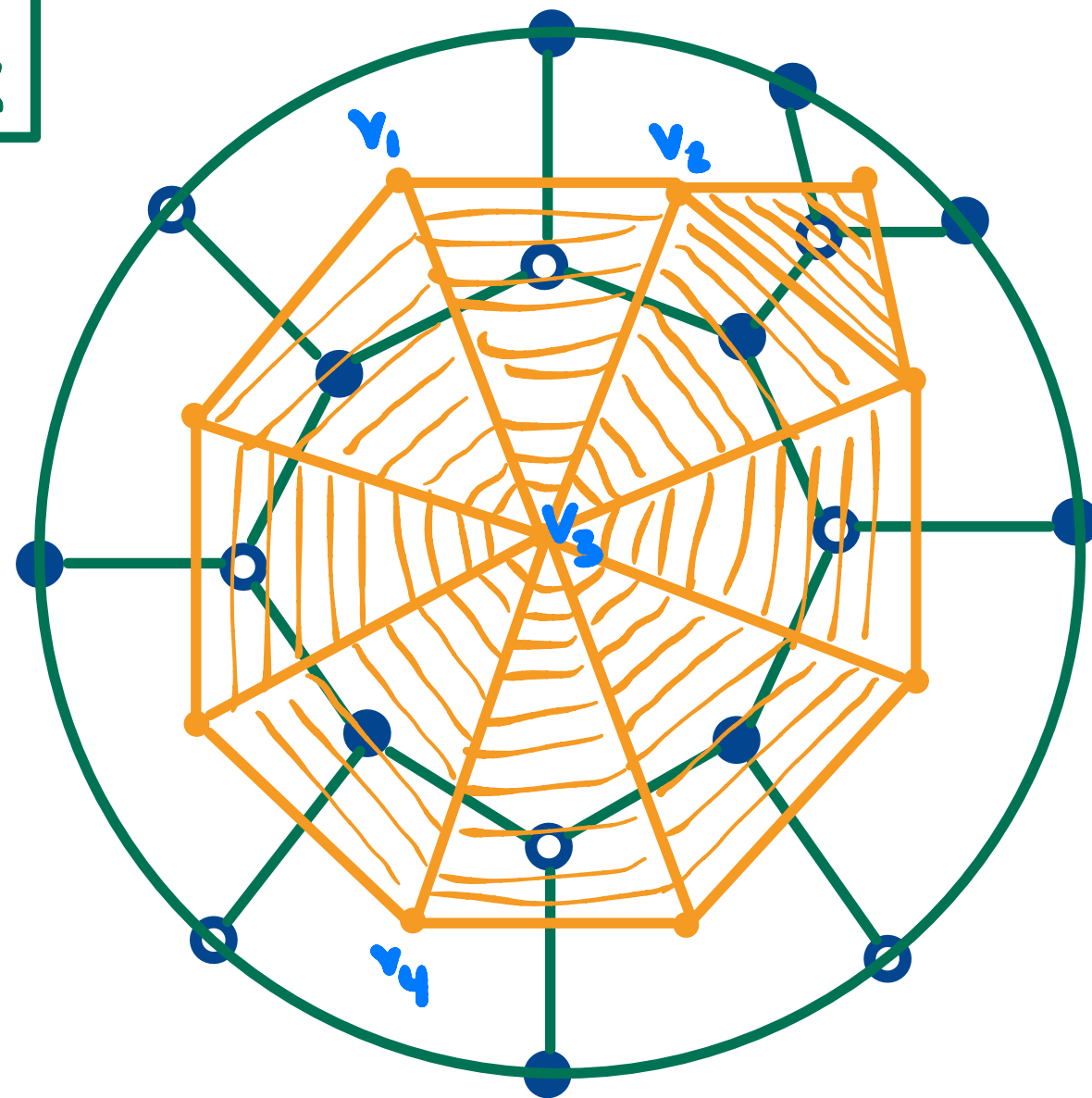
Distances

Idea for an $SL(3)$ web, walk along the dual graph, picking up $SL(3)$ weights according to:



Call this the **distance** of the walk.

Ex



$$d(v_1 \xrightarrow{w_1} v_2 \xrightarrow{w_1} v_3 \xrightarrow{w_2} v_4) = 2w_1 + w_2 = (3, 1, 0)$$

vs.

$$d(v_1 \xrightarrow{w_2} v_3 \xrightarrow{w_2} v_4) = 2w_2 = (2, 2, 0)$$

smaller!

Geodesics

Def • Such a walk on $D(W)$ for an $SL(3)$ -web W is a ("combinatorial") **geodesic** from p to q if it is a minimizer of

$$\{d(\gamma) \mid \gamma: p \rightarrow \dots \rightarrow q\}.$$

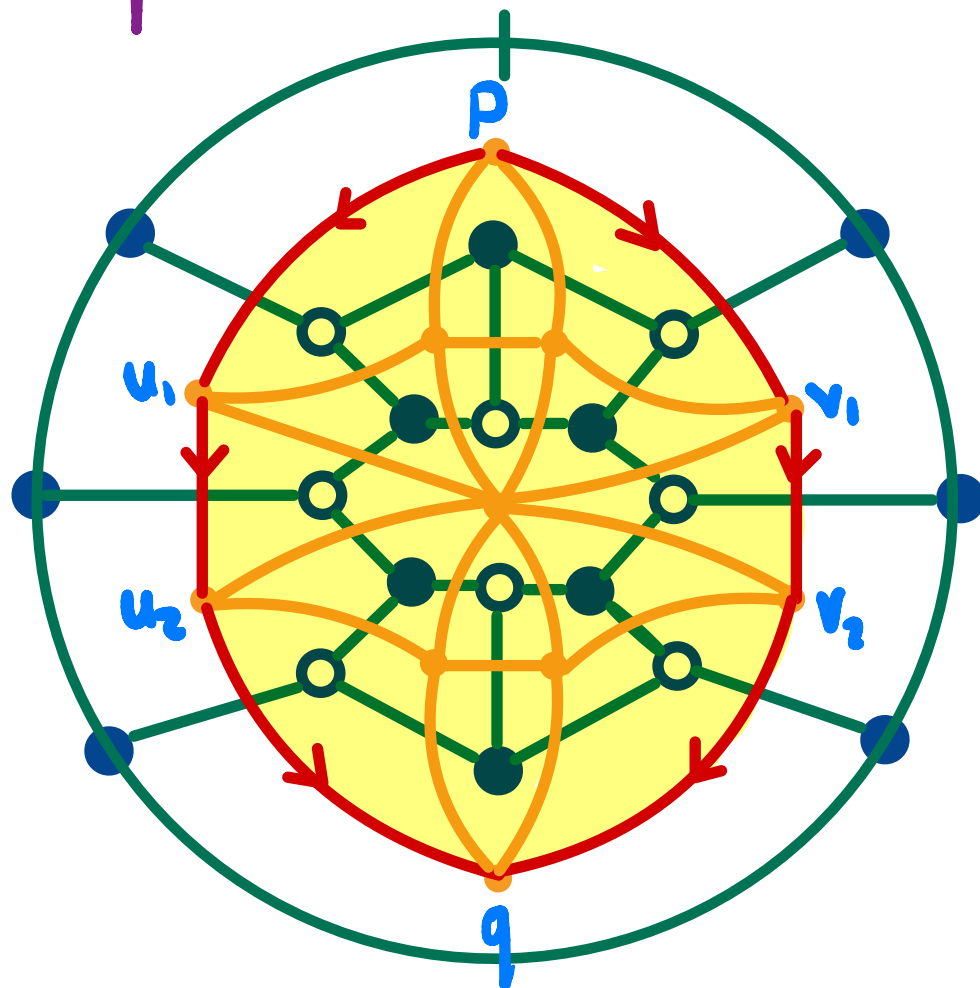
• We have **coherent geodesics** if

$$\forall p, q, \underbrace{\exists!}_{\text{unique!}} \text{ geodesic from } p \text{ to } q.$$

Geodesics

Thm (FKK) IF W is a non-elliptic $SL(3)$ web, it has coherent geodesics.

Ex Non-example:



$$d(p \rightarrow u_1 \rightarrow u_2 \rightarrow q) = (3, 3, 0)$$

$$d(p \rightarrow v_1 \rightarrow v_2 \rightarrow q) = (3, 0, 0)$$

$$= (4, 1, 1)$$

incomparable!

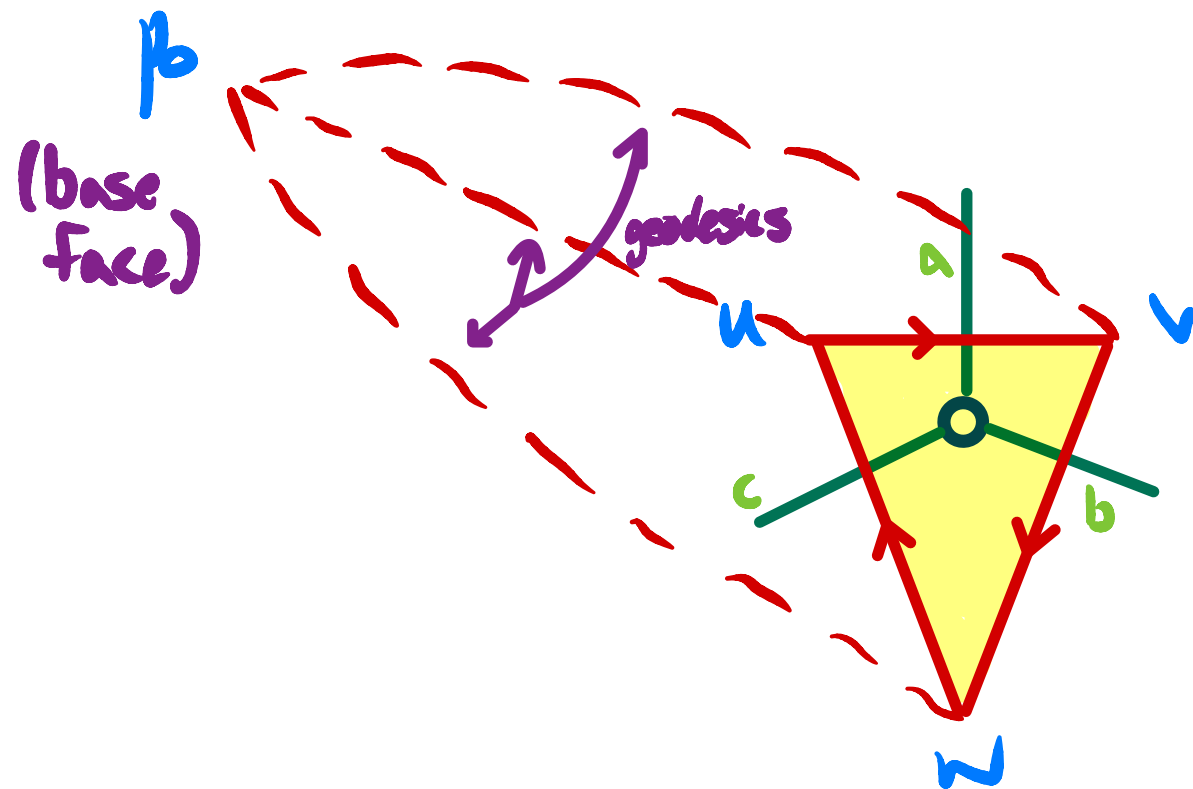
(has squares)

Geodesics

Fact | Geodesic distances from base face to faces around the boundary encodes original $T \in \text{SVT}(3 \times \frac{n}{3})!$

Q | Growth rule label meaning?

A



Fact | $d(p, w) - d(p, v) \in S_r \cdot w$

Hence $d(p, w) - d(p, v) = e_a$
 $d(p, v) - d(p, w) = e_b$
 $d(p, w) - d(p, u) = e_c$

$$0 = e_a + e_b + e_c$$

$$\Leftrightarrow \{a, b, c\} = \{1, 2, 3\}$$

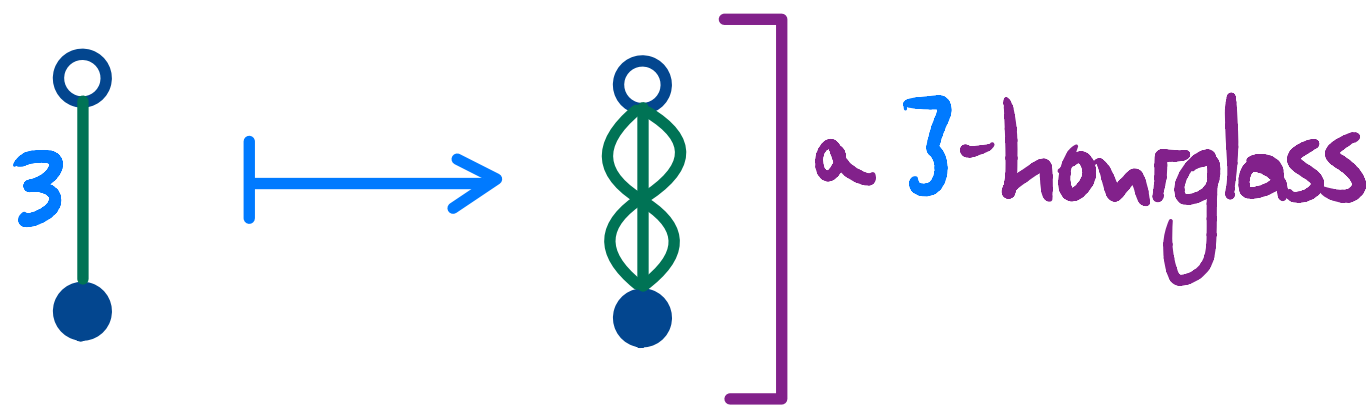
$\Leftrightarrow$ proper labeling!

Hourglass plabic graphs

Def ([Gaetz-Pechenik-Pfannerer-Striker-S. '23])

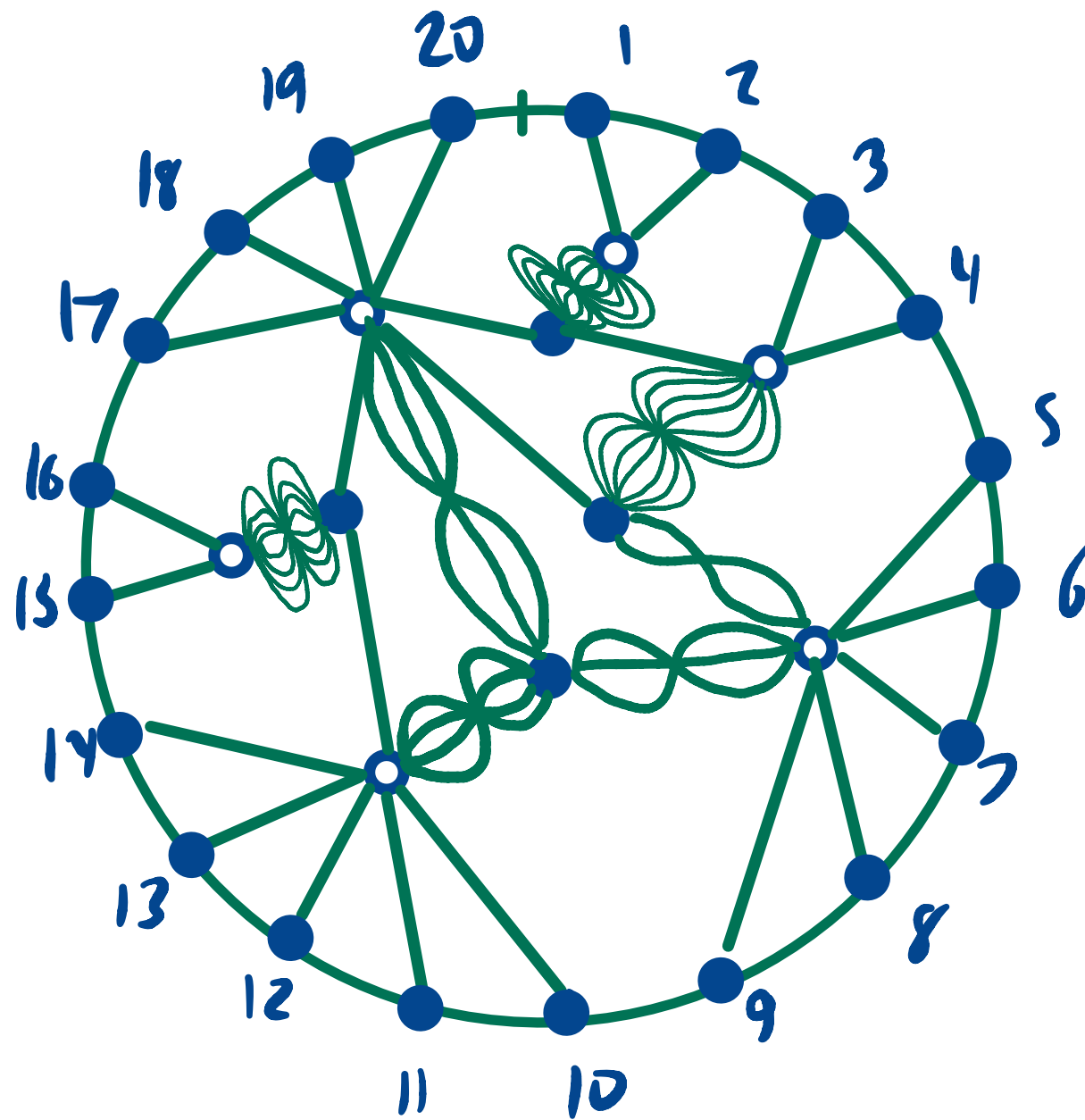
An r -hourglass plabic graph (r -HPG) is a planar bipartite graph embedded in a disk with edge weights in $[r]$ which sum to r around internal vertices, and boundary vertices have degree 1.

An edge with weight m is drawn as an m -hourglass:



Hourglass planar graphs

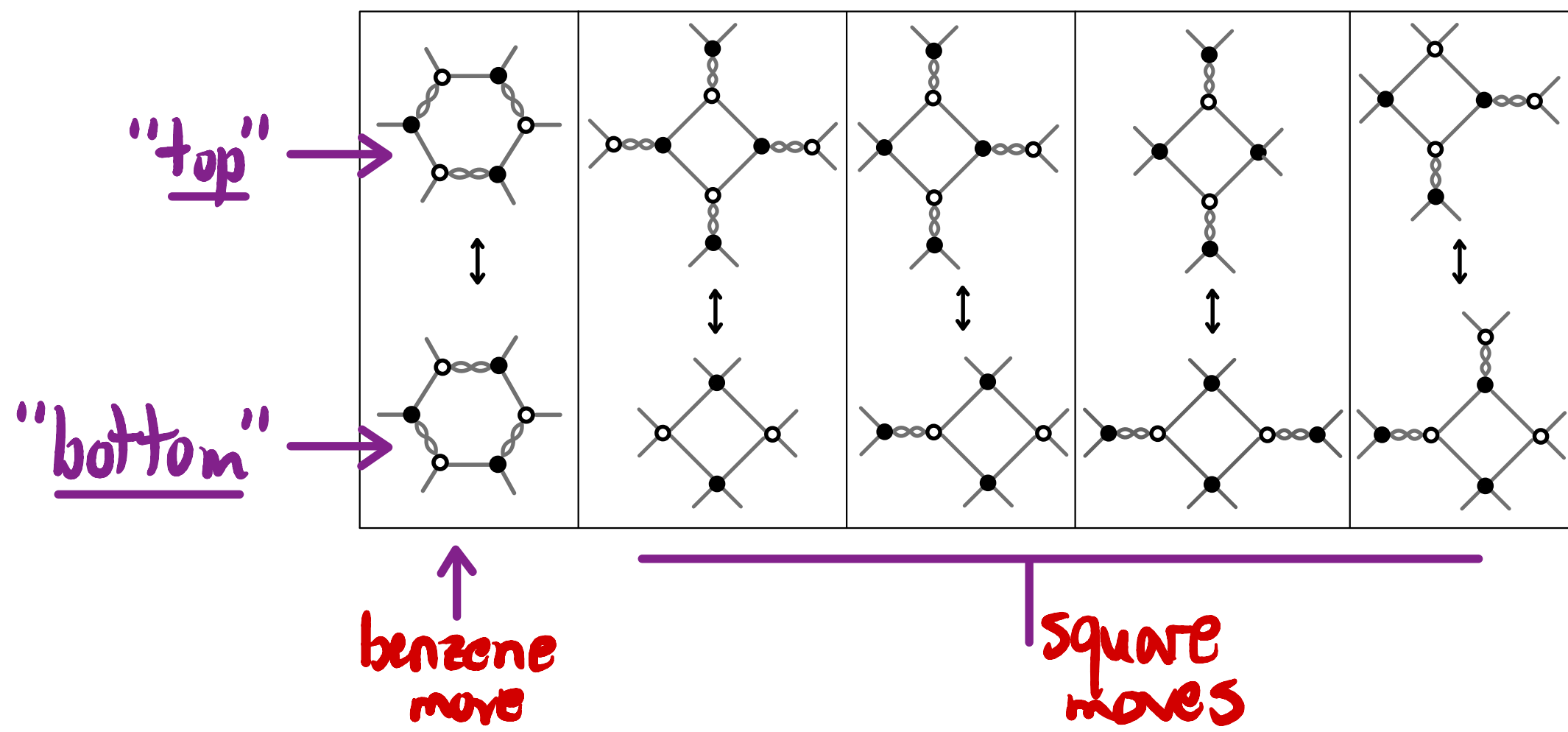
Ex



$(r=10)$

$r=4$ moves

Thm (GPPSS '23) Two contracted, fully reduced
4-HPG's have the same trip permutations
 $\Leftrightarrow$ they are related by moves:



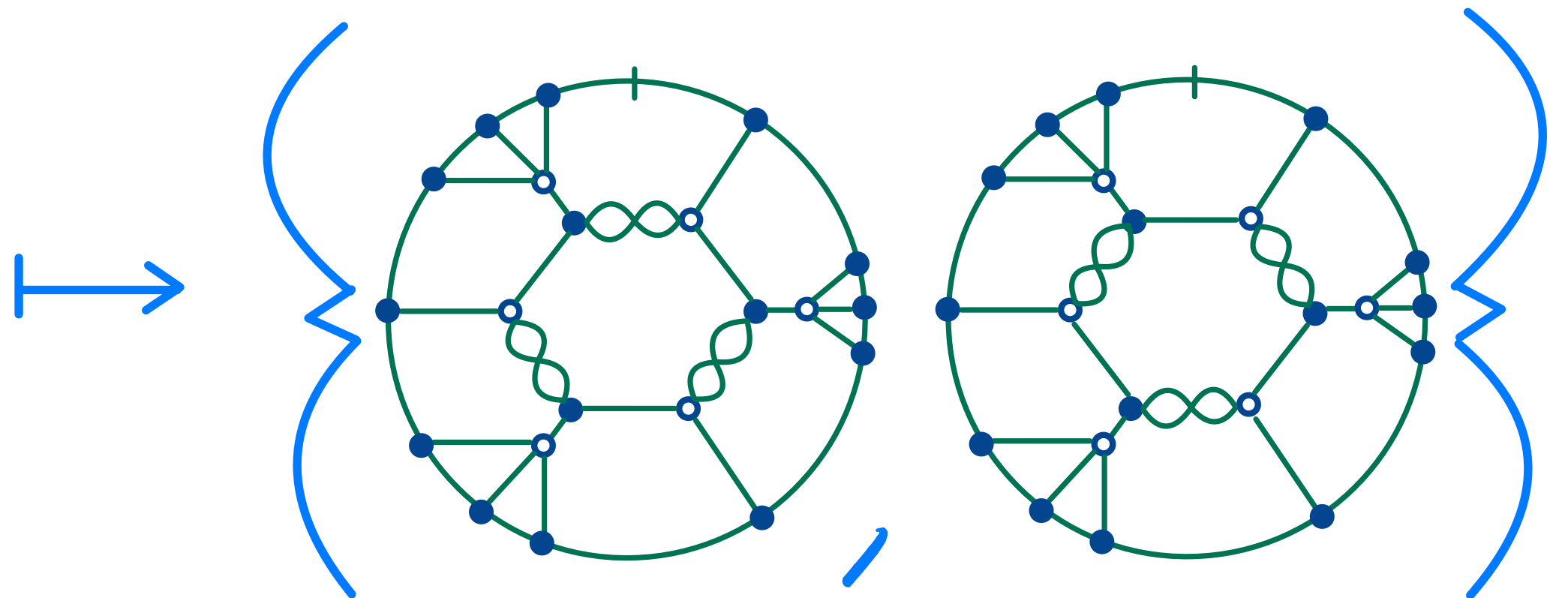
$r=4$ moves

Thm (GPPSS '23) There is a bijection
between $\text{SYT}(4 \times \infty)$ and such move-classes.
It sends $\text{prom}_i(T)$ to $\text{trip}_i(G)$.

Ex

$T =$

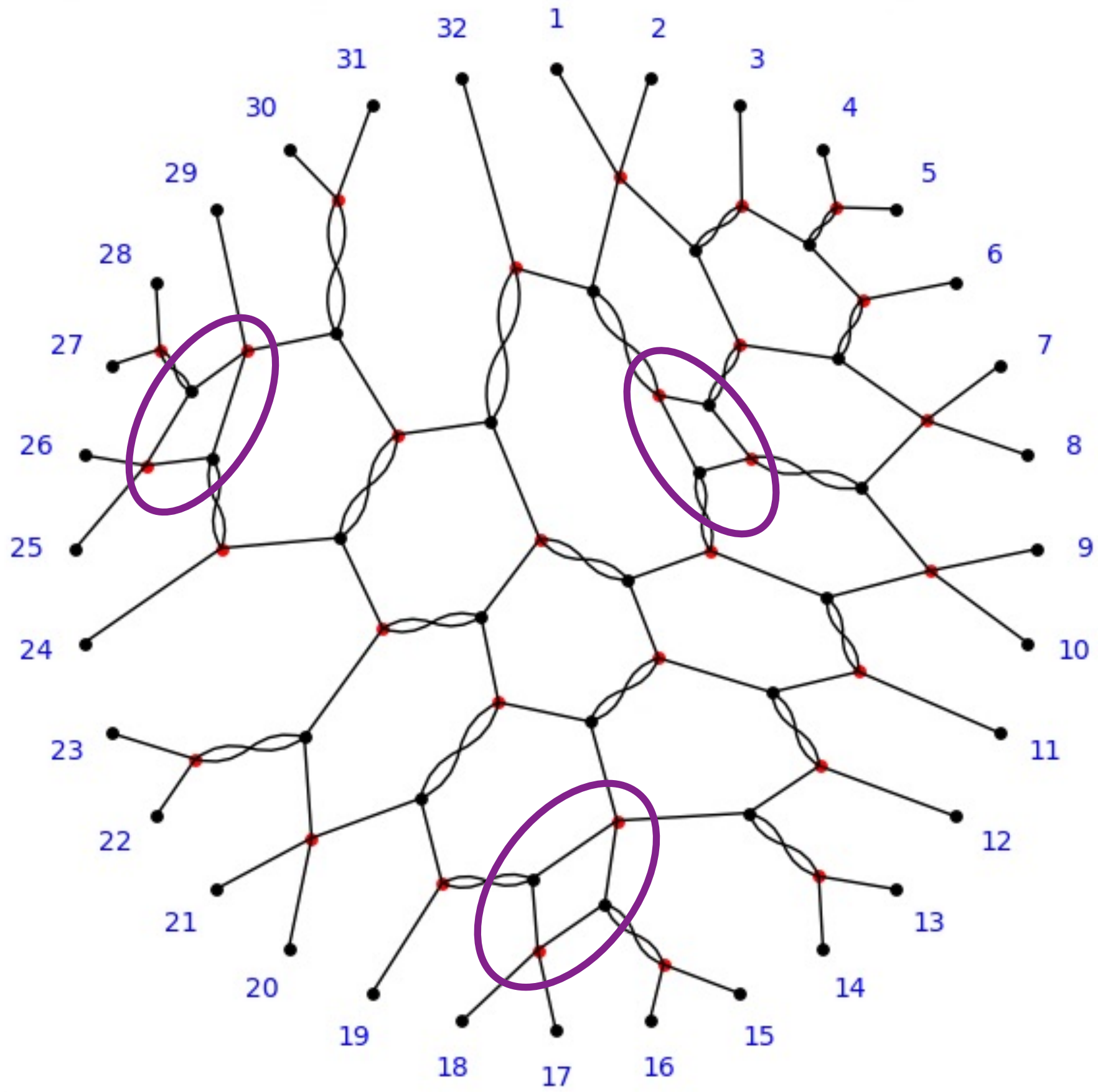
1	2	6
3	5	10
4	7	11
8	9	12



$r=4$ moves

Ex

1	3	4	7	8	17	19	23
2	5	6	9	14	18	21	24
10	12	13	15	16	25	26	28
11	20	22	27	29	30	31	32

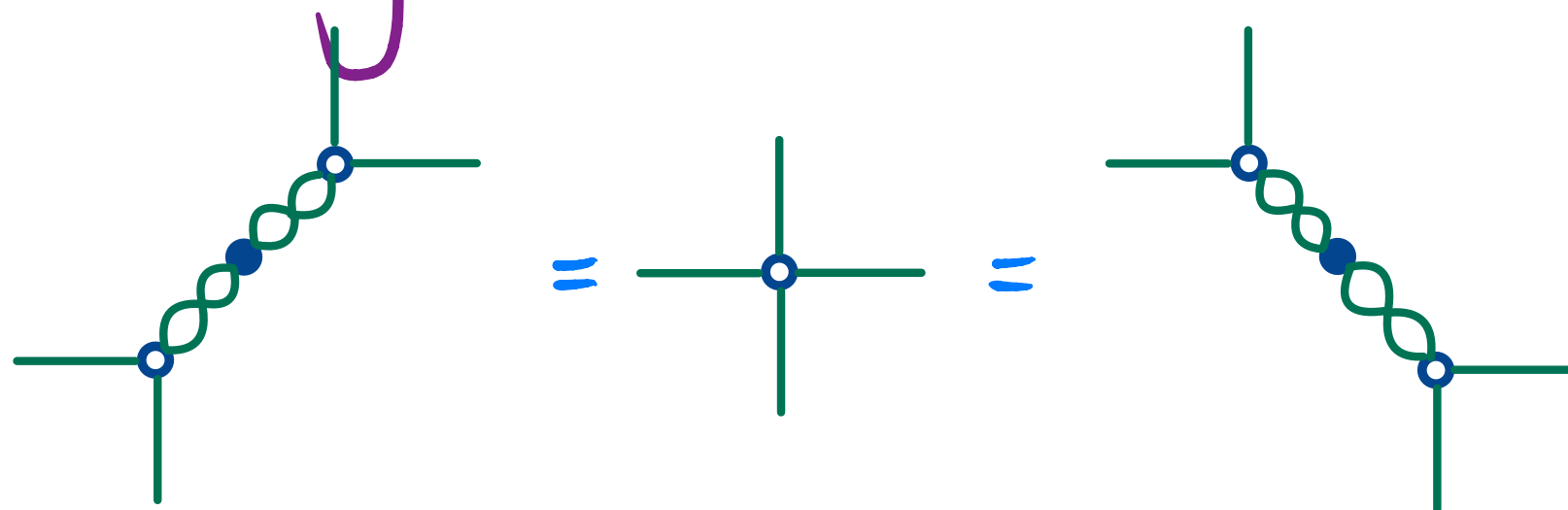


Pockets

Def (Gaetz-Striker-S.-Wu)

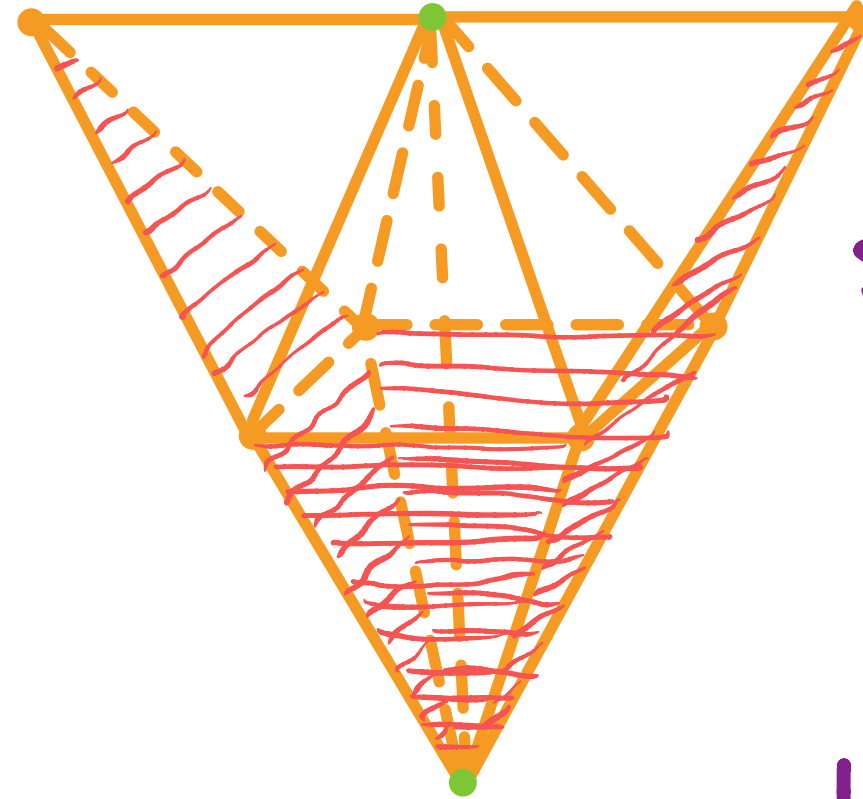
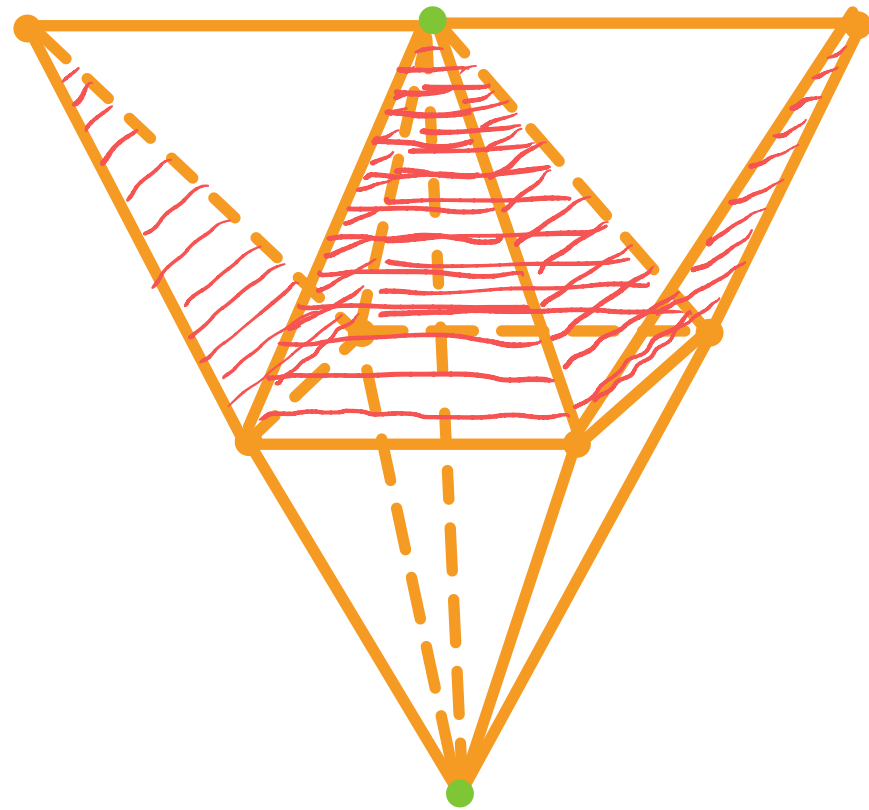
Let $\mathcal{C}$ be the move-class of some fully reduced 4-HPG W . Then there is a 3D flag simplicial complex $P(\mathcal{C})$ called the **pocket** of $\mathcal{C}$, which $D(W)$ embeds into.

- Requires expanding sources/sinks via **I=H moves**:



Pockets

Ex



Pocket of
square + benzene + IH
class of

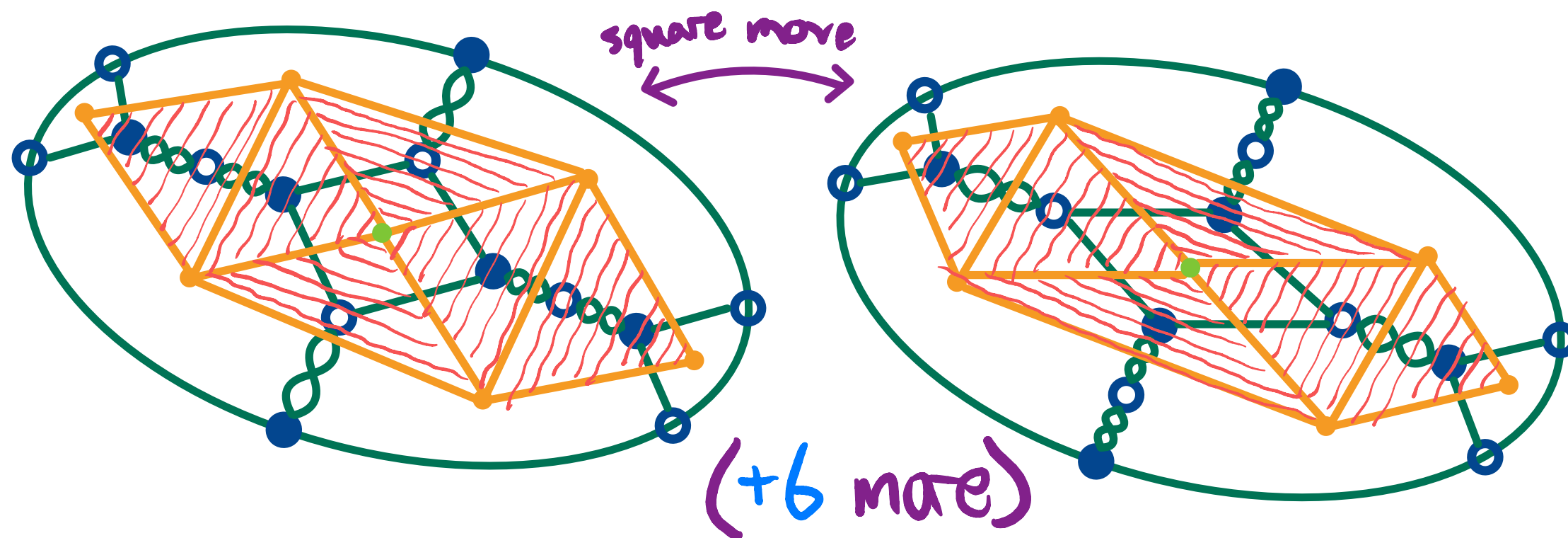
$\overline{4}\{1,2\}\overline{3}\overline{2}\{3,4\}\overline{1}$

has 6 tetrahedra

14 triangles

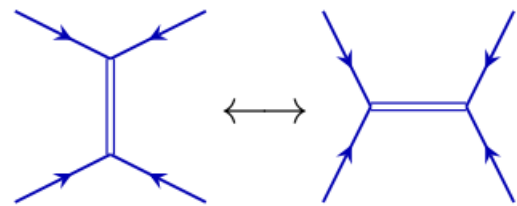
19 edges

8 vertices

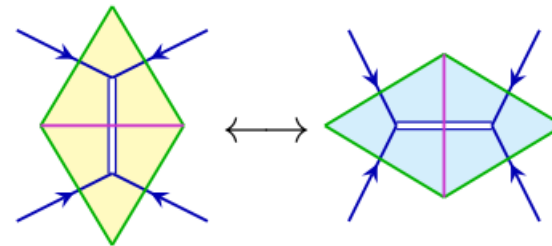


Pochets

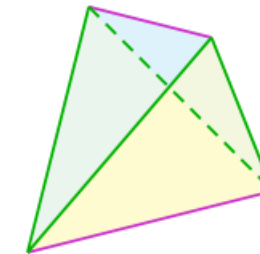
- Build pochets from 4-HPG basis web move classes:



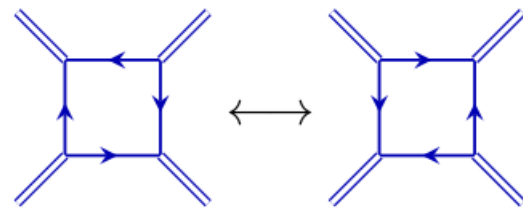
IH move between webs



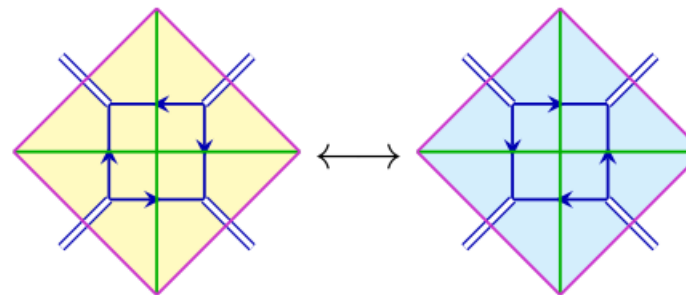
Dual diskoids



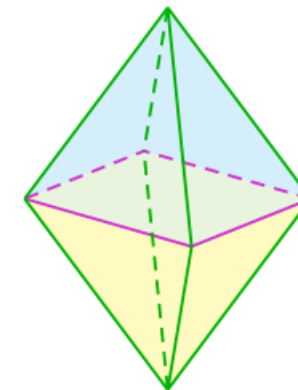
tetrahedron



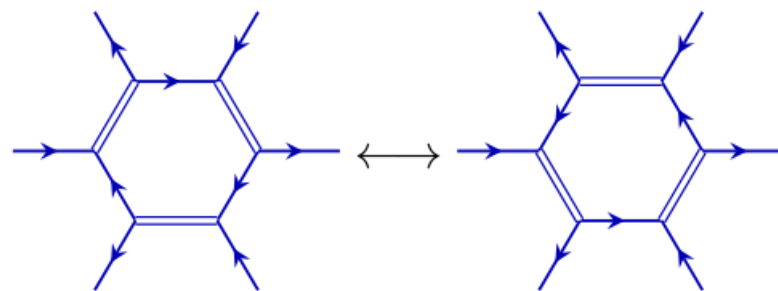
Square move between webs



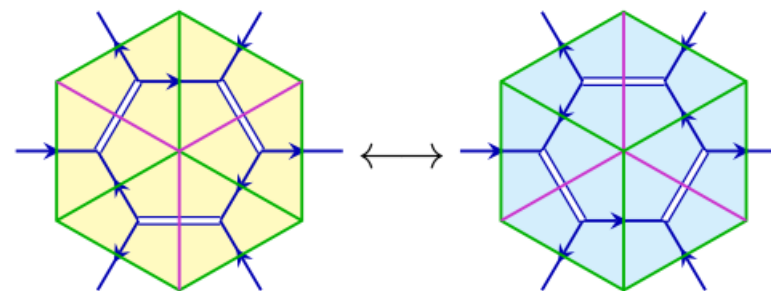
Dual diskoids



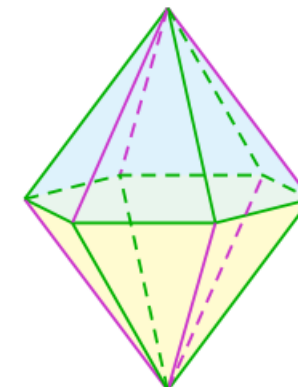
octahedron



Benzene move between webs



Dual diskoids



dodecahedron

Pods

Thm (GSSW)

- The podet $P(\mathcal{C})$ has coherent geodesics.
- It is an interval bundle over the disk.
- The webs $W \in \mathcal{C}$ are its "sections."
- The separation labeling/growth labeling encodes geodesic distances from the base face.

Buildings

Thm (FKK) IF W is a non-elliptic $SL(3)$ web, then $D(W)$ embeds in the affine building $\Delta(SL(3)^V)$, preserving distances.
next!

— In fact, $D(W)$ is CAT(0).

Double cosets

Obs | Given groups $H \leq G$, have a "distance" on G :
(or G/H)

$$d: G \times G \rightarrow H \backslash G / H$$

$$d(p, q) = \underbrace{Hp^{-1}qH}_{\text{a double coset}}$$

a double coset

Ex | For $\mathbb{R}_{>0} \subseteq \mathbb{C}^*$, (double) cosets are represented by polar angles. Then $d(re^{i\theta}, se^{i\phi}) = e^{i(\phi-\theta)} \mathbb{R}_{>0}$.

Basically $(x, y) \mapsto \arg\left(\frac{y}{x}\right) = \text{atan2}(y, x)!$

Double cosets

Ex For $\mathbb{C}[[t]]^\times < \mathbb{C}(t)$, coset reps are $\{t^n \mid n \in \mathbb{Z}\}$

Ex For $GL_r(\mathbb{C}[[t]]) < GL_r(\mathbb{C}(t))$, double coset reps are

$$t^\lambda = \begin{pmatrix} t^{\lambda_1} & & \\ & t^{\lambda_2} & \\ & & \ddots \\ & & & t^{\lambda_r} \end{pmatrix} \quad \text{where } \lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_r \quad \left. \begin{array}{l} \text{integers.} \\ \text{GL}(r) \\ \text{dominant} \\ \text{weights} \end{array} \right]$$

Ex For $PGL_r(\mathbb{C}[[t]]) < PGL_r(\mathbb{C}(t))$, now have

$$t^\lambda = t^{\lambda + (1^r)} \quad \left. \begin{array}{l} \text{SL}(r) \text{ dominant} \\ \text{weights} \end{array} \right]$$

Affine Grassmannians

Def The affine Grassmannian of $SL(r)^\vee$ is

$$\text{AffGr}_r = \text{PGL}_r(\mathbb{C}(t)) / \text{PGL}_r(\mathbb{C}[t]).$$

- Have "distance"

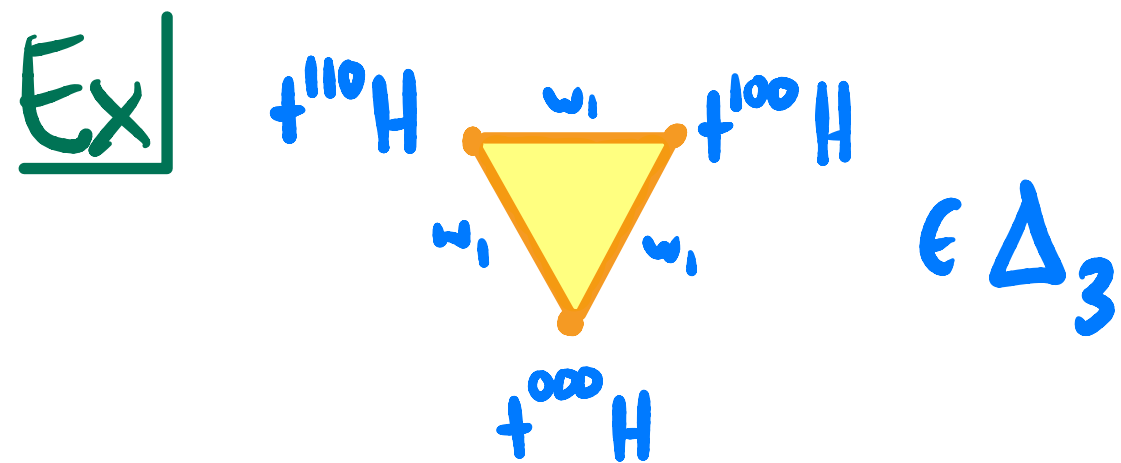
$$d: \text{AffGr}_r \times \text{AffGr}_r \rightarrow \mathbb{Z}_{\text{dec}}^r / \langle 1, \dots, 1 \rangle \Bigg] \begin{array}{l} \text{SL}(r) \text{ dominant} \\ \text{weights} \end{array}$$

$$\text{where } d(pH, qH) = d(p, q) = \lambda \Leftrightarrow p^{-1}q \in Ht^\lambda H$$

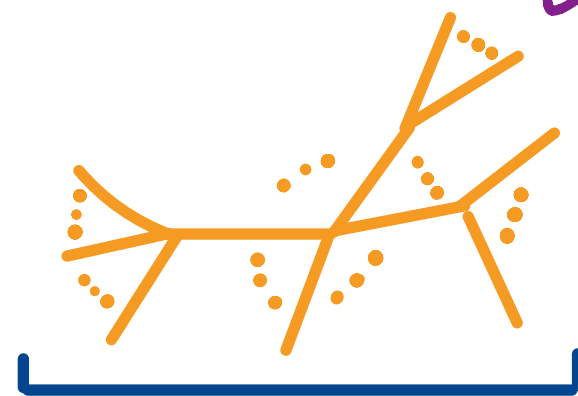
- Have $d(p, p) = 0$ $d(p, q) = d(gp, gq)$ $d(q, p) = \underline{-\text{rev}(d(p, q))}$
 for $g \in \text{PGL}_r(\mathbb{C}(t))$ not quite (anti)symmetric

Affine Buildings

Def The affine building on AffGr_r is the simplicial complex Δ_r whose vertices are the points of AffGr_r and whose simplices are collections of points all of whose distances are fundamental weights.



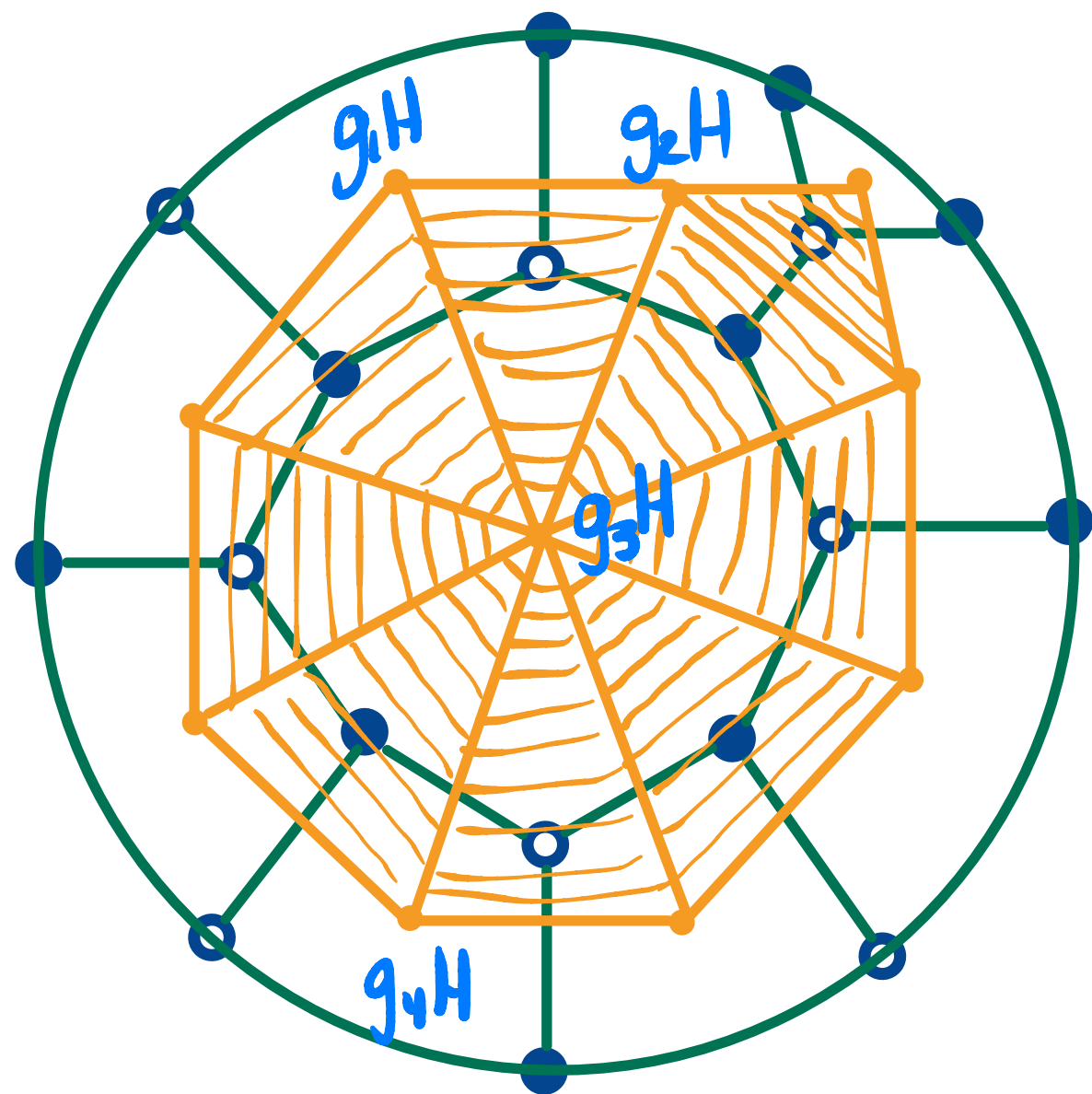
Ex $\Delta_2 =$



tree where every vertex has uncountable degree

Affine Buildings

Ex | There exist $g_1H, g_2H, g_3H, g_4H, \dots \in \text{AffGr}$, s.t.



$$d(g_1, g_2) = w_1$$

$$d(g_2, g_3) = w_1$$

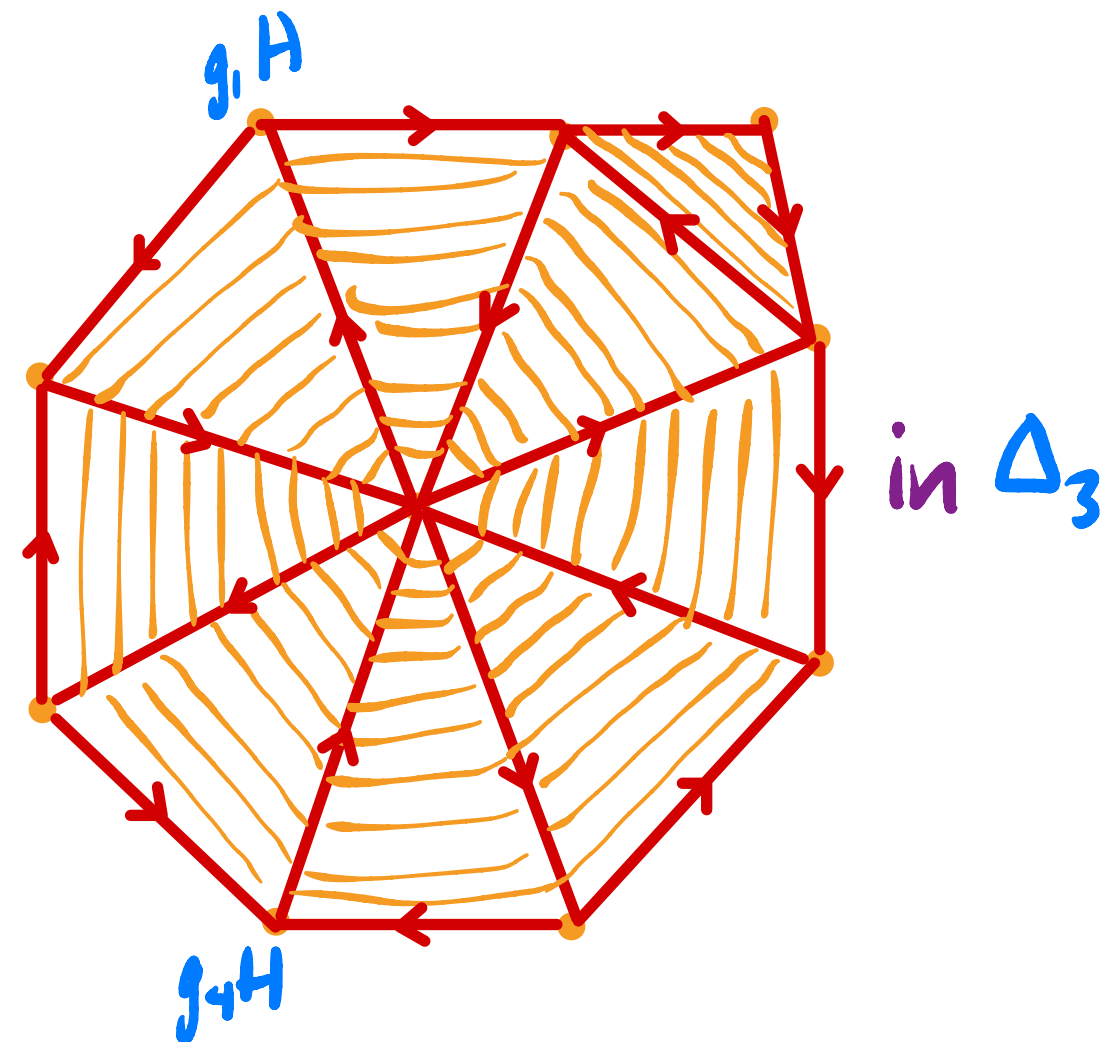
$\vdots$

$$d(g_4, g_1) = \underline{2w_1}$$

$\vdots$

geodesic
distance

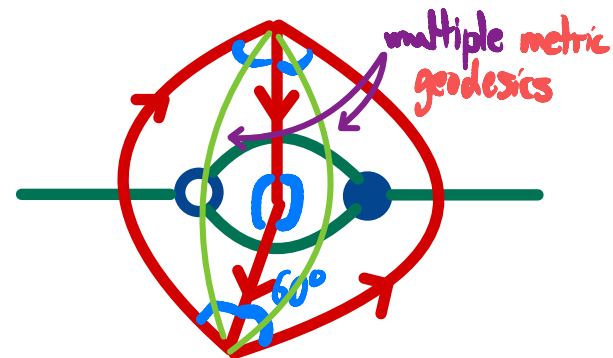
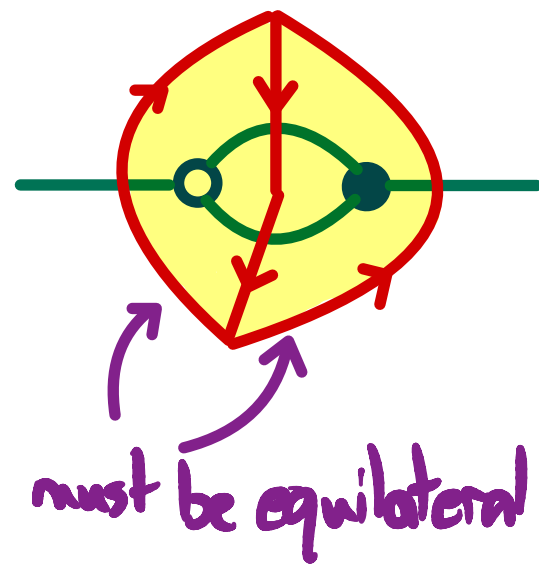
$$\begin{array}{l} \xrightarrow{\text{red}} = w_1 \\ \xleftarrow{\text{purple}} = w_1^* = w_2 \end{array}$$



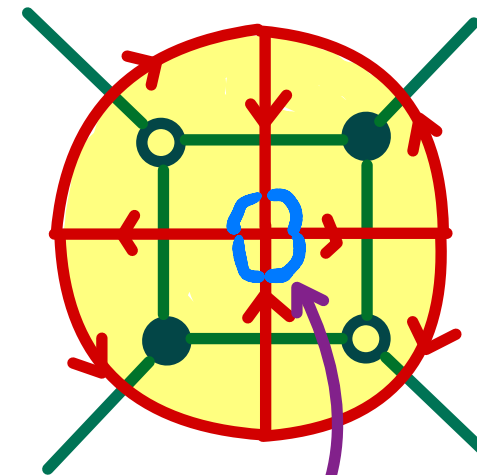
Affine Buildings

Q Non-elliptic condition meaning?

A



Not CAT(0)!



Angle $4 \cdot 60^\circ = 240^\circ < 360^\circ$!
Also not CAT(0)!

Pockets

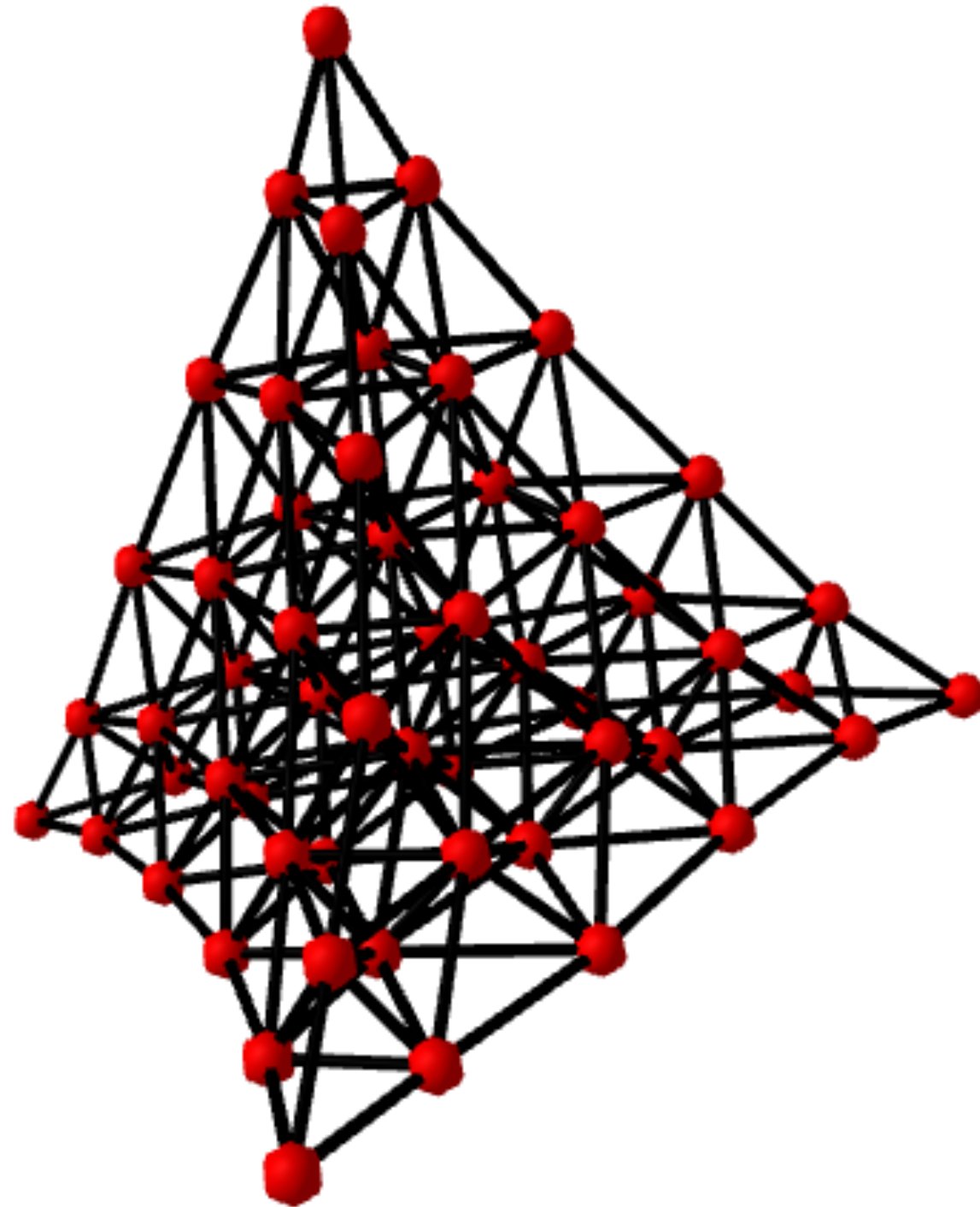
Thm (GSSW)

- Pockets $P(\mathcal{C})$ embed in Δ_4 .
- Indeed, $P(\mathcal{C})$ is CAT(0).
- Given an irreducible component of a Satake fiber of Δ_4 indexed by T , there is a dense open subset U s.t. every point of U extends uniquely to a configuration $P(\mathcal{C}) \hookrightarrow \Delta_4$ which preserves distances.

Podnets

Ex Product of 5×5 ASM:

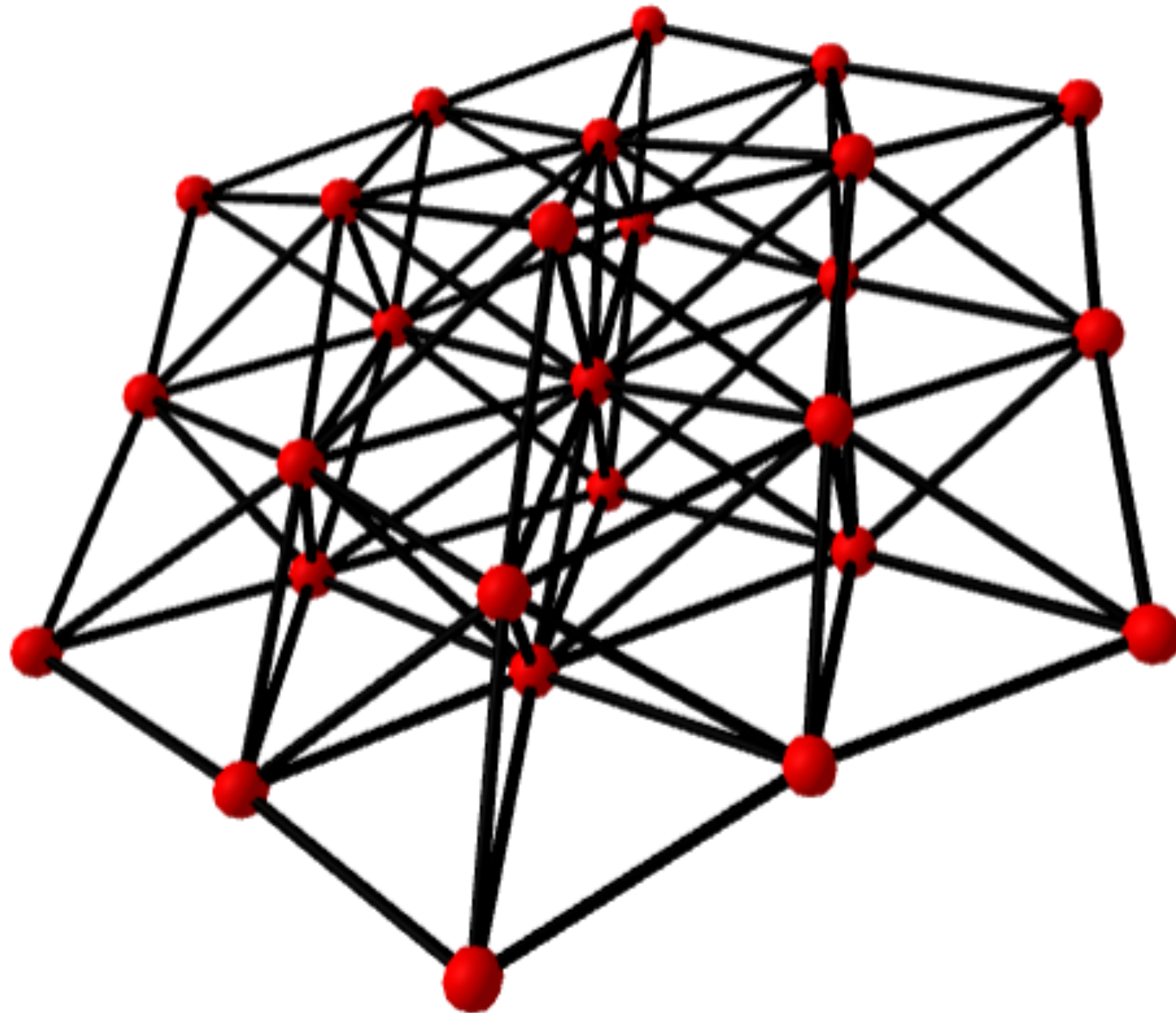
$$L = 1^5 2^5 3^5 4^5$$



Note Related to
height functions,
octahedral
recurrence,
tilings of the
Aztec diamond,
distributive lattices

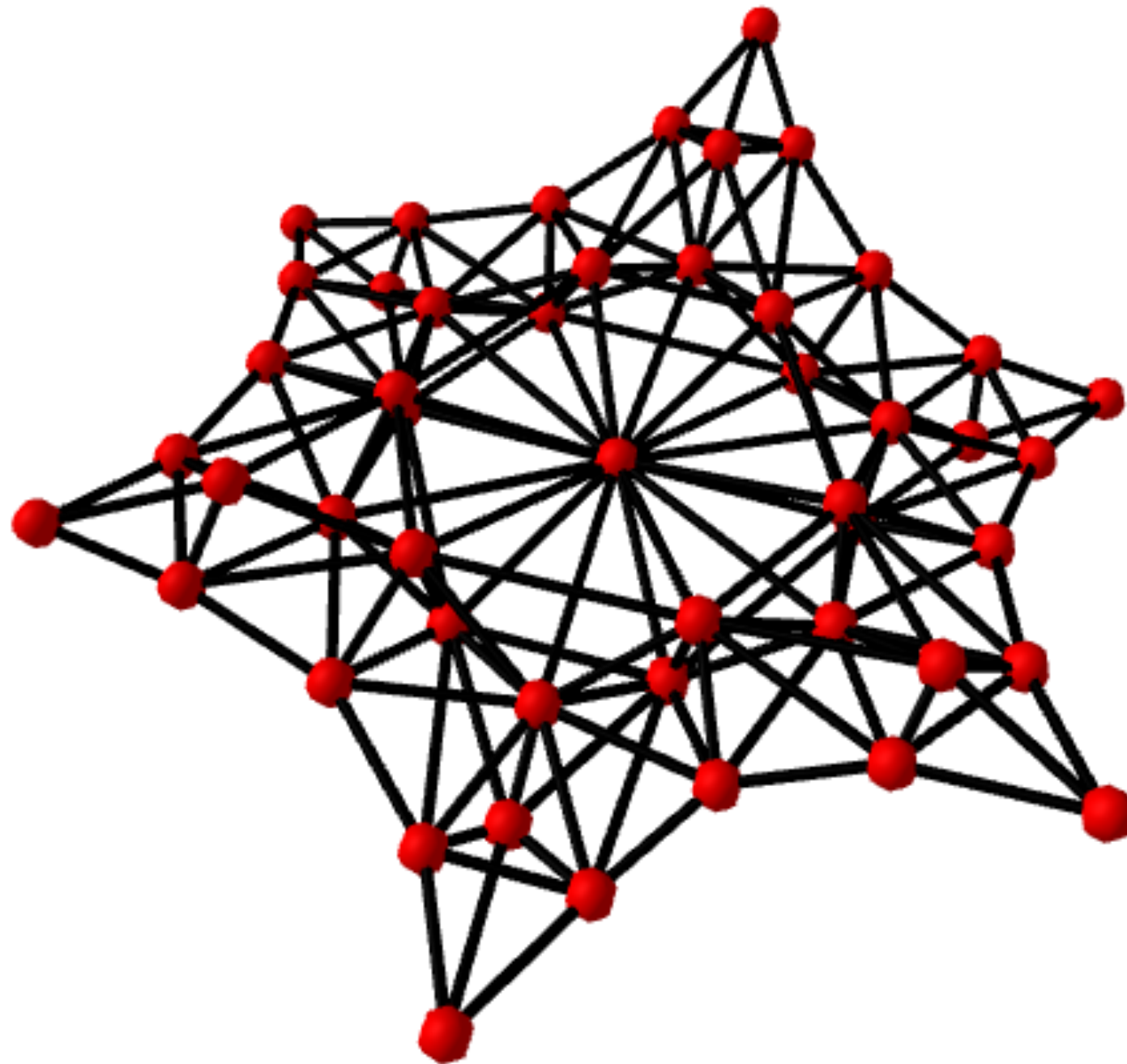
Products

Ex Product of $2 \times 2 \times 2$ PP:



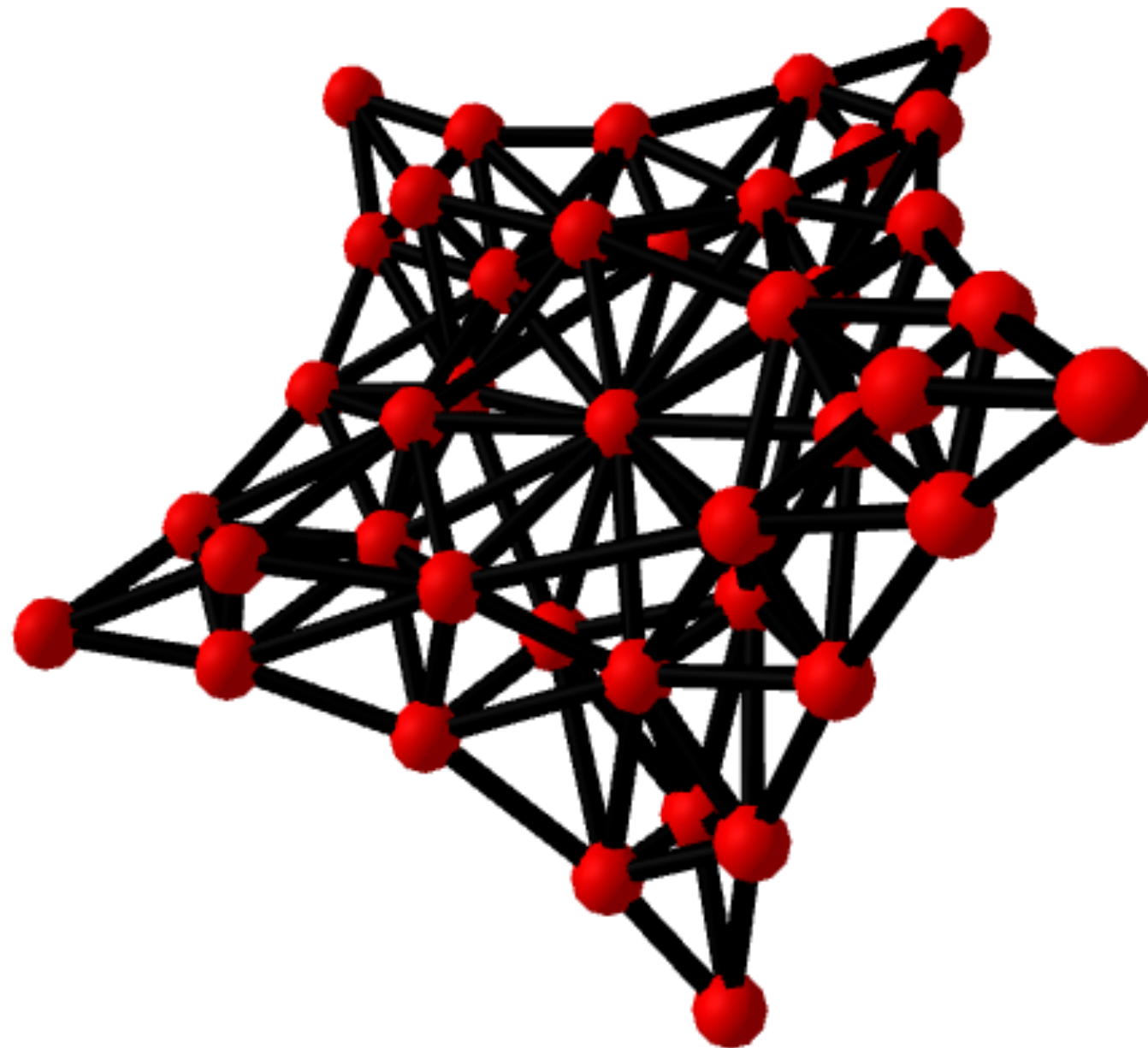
Podsets

Ex Product of a "chained hexagon":



Podnets

Ex | Product of a "chained pentagon":



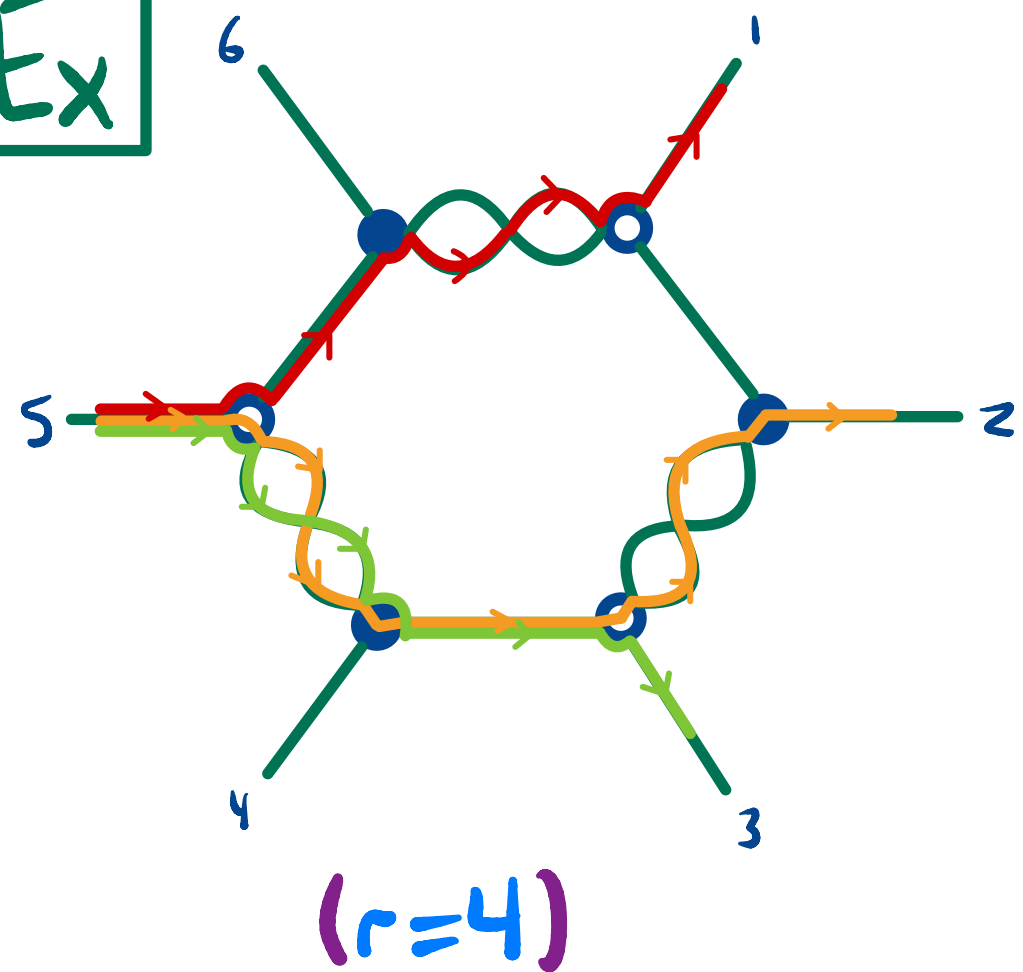
Note | Not realizable
in $\mathbb{R}^3$

THANKS!

Trip permutations

Def (GPPSS '23) An r -hourglass planar graph has trip permutations $\text{trip}_1, \dots, \text{trip}_{r-1}$ where trip_i takes the i th left at white and i th right at black:

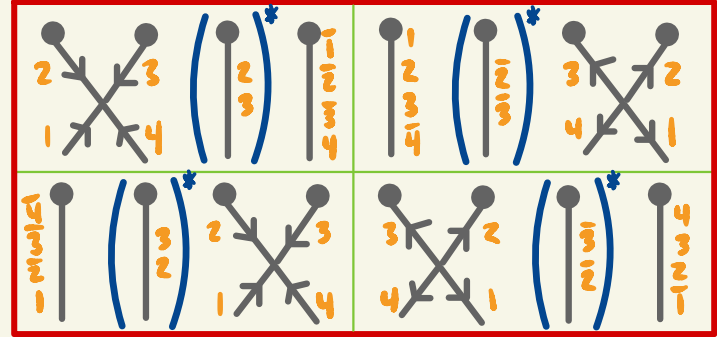
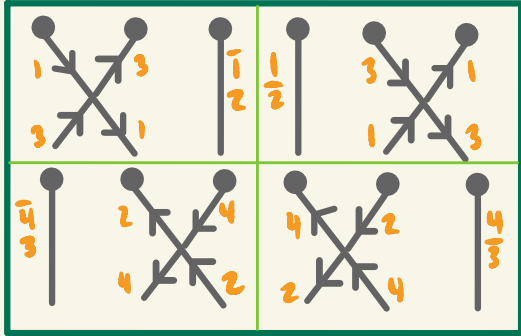
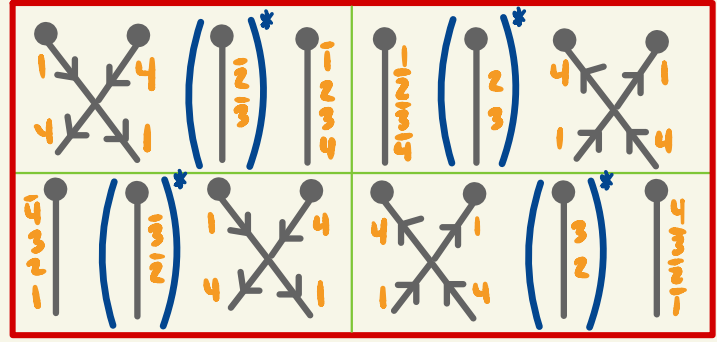
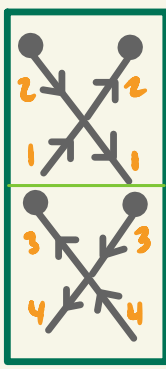
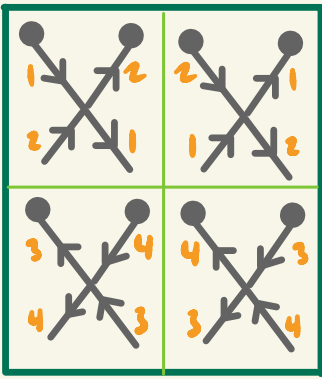
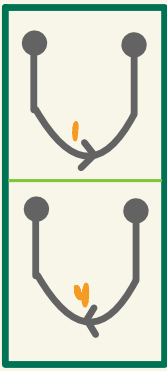
Ex



$\rightarrow$ = $\text{trip}_1 = (1\ 3\ 5)(6\ 4\ 2)$
 $\rightarrow$ = $\text{trip}_2 = (1\ 4)(2\ 5)(3\ 6)$
 $\rightarrow$ = $\text{trip}_3 = (5\ 3\ 1)(2\ 4\ 6)$

Note

$\text{trip}_i = \text{trip}_{r-i}^{-1}$!



$SL(4)$ growth rules

